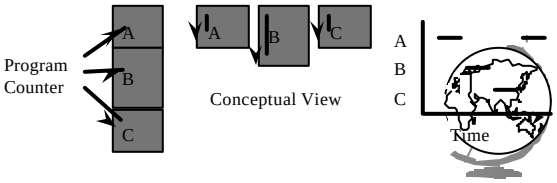
Operating Systems

Processes
(Ch 4.1)

Processes

- "A program in execution"
- Modern computers allow several at once
  - "pseudoparallelism"



Processes

- "A program in execution"

```
main() {
...
}
A() {
…
}
```

```
main() {
...
}
A() {
…
}
```

Heap

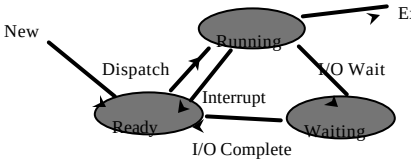
Stack

A main

  - "more" than a program: `ls, tcsh`
  - "less" than a program: `gcc blah.c`
    `(cpp, cc1, cc2, ln …)`
- "A sequential stream of execution in it's own address space"

Process States

- Consider:
  `cat /etc/passwd | grep claypool`



(Hey, you, show states in `top`!)

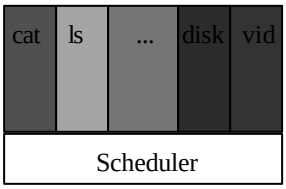
Design Technique: State Machines

- Process states
- Move from state to state based on events
  - *Reactive* system
- Can be mechanically converted into a program
- Other example:
  - string parsing, pre-processor

Unix Process Creation

- System call: `fork()`
  - creates (nearly) identical copy of process
  - return value different for child/parent
- System call: `exec()`
  - over-write with new process memory
- Shell
  - uses fork() and exec()
  - simple!
- (Hey, you, show demos!)

Process Scheduler



- All services are processes
- Small scheduler handles interrupts, stopping and starting processes

Process Control Block

- Each process has a PCB
  - state
  - program counter
  - registers
  - memory management
  - …
- OS keeps a table of PCB's, one per process
- (Hey! Simple Operating System, "`system.h`")

Question

- Usually the PCB is in OS memory only.
- Assume we put the PCB into a processes address space.
- What problems might this cause?

Interrupt Handling

- Stores program counter (hardware)
- Loads new program counter (hardware)
  - jump to interrupt service procedure
- Save PCB information (assembly)
- Set up new stack (assembly)
- Set "*waiting*" process to "*ready*" (C)
- Re-schedule (probably awakened process) (C)
- If new process, called a *context-switch*

Context Switch

- Pure overhead
- So … fast, fast, fast
  - typically 1 to 1000 microseconds
- Sometimes special hardware to speed up
- How to decide when to switch context to another process is *process scheduling*

Processes in Linux

- PCB is in `struct task_struct`
  - states: RUNNING, INTERRUPTIBLE, UNINTERRUPTIBLE
  - priority: when it runs
  - counter: how long it runs
- Environment inherited from parent
- NR_TASKS max, 2048
  - 1/2 is max per user

Processes in NT

- States: ready, standby (first in line), running, waiting, transition, terminated
- priority - when it runs
- Processes are composed of *threads*
  - (revisit threads after scheduling)

True or False

- Unix is a “simple structure” OS
- Micro Kernels are faster than other OS structures
- Virtual Machines are faster than other OS structures