

Operating System I

Process Scheduling

Schedulers

- **♦** Short-Term
 - "Which process gets the CPU?"
 - Fast, since once per 100 ms
- **→** Long-Term (batch)
 - "Which process gets the Ready Queue?"
- → Medium-Term (Unix)
 - "Which Ready Queue process to memo
 - Swapping



CPU-IO Burst Cycle add read (I/O Wait) store increment write (I/O Wait) Burst Duration

Preemptive Scheduling

- → Four times to re-schedule
 - 1 Running to Waiting (I/O wait)
 - 2 Running to Ready (time slice)
 - 3 Waiting to Ready (I/O completion)
 - 4 Termination
- ◆ #2 and #3 optional ==> "Preemptive"
- → Timing may cause unexpected result
 - updating shared variable
 - kernel saving state



Question

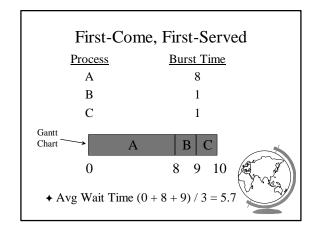
- ♦ What Performance Criteria Should the Scheduler Seek to Optimize?
 - Ex: CPU minimize time spent in queue
 - Others?

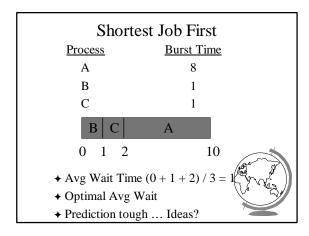


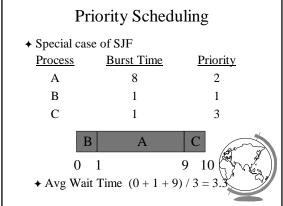
Scheduling Criteria

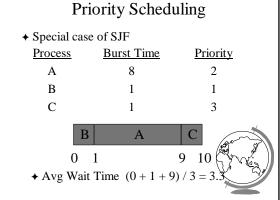
- 1 CPU utilization (40 to 90)
- 2 Throughput (processes / hour)
- 3 Turn-around time
- 4 Waiting time (in queue)
- ◆ Maximize #1, #2 Minimize #3, #4
- **→** Response time
 - Self-regulated by users (go home)
 - Bounded ==> Variance!











- \$ - process importance - <u>super-user (root)</u> - <u>nice</u> Round Robin SOS: Dispatcher ◆ Fixed time-slice and Preemption **Process Burst Time** Α 5

В 3 \mathbf{C} В Avg = (8 + 9 + 11) / 3 = 9.3◆ FCFS? SJF?

Priority Scheduling Criteria? **→** Internal - open files - memory requirements - CPU time used - time slice expired (RR) - I/O wait completed - process age ◆ External - department sponsoring work

→ How is the next process chosen? → Line 79 has an infinite loop. Why? ◆ There is no return from the Dispatcher() function call. Why not? ◆ See "TimerInterruptHandler()" - /usr/src/linux/kernel/sched.c - /usr/src/linux/include/linux/sd - linux-pcb.h

Round Robin Fun

<u>Process</u>	Burst Time
A	10
В	10
C	10

- **→** Turn-around time?
 - q = 10
 - q = 1
 - -q > 0



More Round Robin Fun Process Burst Time A 6 B 3 C 1 D 7 Rule: 80% within one quantum Rule: 1 2 3 4 5 6 7 Time Quantum

Fun with Scheduling

<u>Process</u>	Burst Time	Priority
A	10	2
В	1	1
C	2	3

- **→** Gantt Charts:
 - FCFS
 - SJF
 - Priority
 - -RR(q=1)

◆ Performance: – Throughput – Waiting time – Turnaround time

More Fun with Scheduling

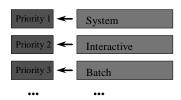
Process	Arrival Time	Burst Time
A	0.0	8
В	0.4	4
C	1.0	1

- **→** Turn around time:
 - FCFS
 - SJF
 - q=1 CPU idle
 - q=0.5 CPU idle



Multi-Level Queues

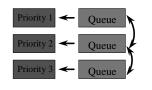
◆ Categories of processes



- → Run all in 1 first, then 2 ...
- **→** Starvation!
- → Divide between queues: 70% 1, 15% 2

Multi-Level Feedback Queues

→ Time slice expensive but want interactive



- 1 Quantum
- 2 Quanta
- 4 Quanta
- ◆ Consider process needing 100 quanta
- -1, 4, 8, 16, 32, 64 = 7 swaps!
- ◆ Favor interactive users



Windows NT Scheduling

- → Basic scheduling unit is a thread
- ◆ Priority based scheduling per thread
- ◆ Preemptive operating system
- → No shortest job first, no quotas



Priority Assignment

- → NT kernel uses 31 priority levels
 - 31 is the highest; 0 is system idle thread
 - Realtime priorities: 16 31
 - Dynamic priorities: 1 15
- ◆ Users specify a priority class:
 - realtime (24), high (13), normal (8) and idle (4)
 - and a relative priority:
 - highest (+2), above normal (+1), normal (0) performance (-1), and lowest (-2)
 - to establish the Starting priority
- → Threads also have a current priority

Quantum

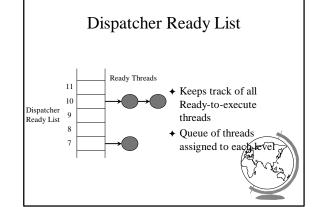
- ◆ Determines how long a Thread runs once selected
- → Varies based on:
 - NT Workstation or NT Server
 - Intel or Alpha hardware
 - Foreground/Background application threads
 - NOTE: NT 4.0 increases quantum for foregoind threads while NT 3.5 increased priorities.



Outline Processes PCB Interrupt Handlers Scheduling Algorithms WinNT Linux

Questions

- → True or False:
 - FCFS is optimal in terms of avg waiting time
 - Most processes are CPU bound
 - The shorter the time quantum, the better
- ♦ What is the *idle thread*? Where did we see it?



FindReadyThread

- ◆ Locates the highest priority thread that is ready to execute
- ◆ Scans dispatcher ready list
- ◆ Picks front thread in highest priority nonempty queue
- ♦ When is this like round robin?



Boosting and Decay

- **♦** Boost priority
 - Event that "wakes" blocked thread
 - Boosts never exceed priority 15 for dynamic
 - Realtime priorities are not boosted
- ◆ Decay priority
 - by one for each quantum
 - decays only to starting priority (no low



Starvation Prevention

- ◆ Low priority threads may never execute
- → "Anti-CPU starvation policy"
 - thread that has not executed for 3 seconds
 - boost priority to 15
 - double quantum
- → Decay is swift not gradual after this boost



Linux Process Scheduling

- → Two classes of processes:
 - Real-Time
 - Normal
- **♦** Real-Time:
 - Always run Real-Time above Normal
 - Round-Robin or FIFO
 - "Soft" not "Hard"



Linux Process Scheduling

- ◆ Normal: Credit-Based
 - process with most credits is selected
 - time-slice then lose a credit (0, then suspend)
 - no runnable process (all suspended), add to every process:

credits = credits/2 + priority

◆ Automatically favors I/O bound proce