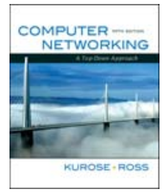
Introduction

CS 3516 – Computer Networks

Chapter 1 Introduction



5th edition
Jim Kurose, Keith Ross
Addison-Wesley, April 2009

Chapter 1: Introduction

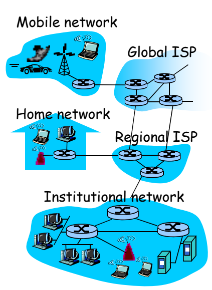
Goal:
- Get "feel" and terminology
- More depth, detail *later* in course
- Approach:
  - use Internet as example

Overview:
- What's the Internet?
- What's a protocol?
- Network edge; hosts, access net, physical media
- Network core: packet/circuit switching, Internet structure
- Performance: loss, delay, throughput
- Security
- Protocol layers, service models
- history

Chapter 1: Roadmap

What's the Internet: "Nuts and Bolts" view

- Millions of connected computing devices: *hosts = end systems*
  - running *network apps*
- *Communication links*
  - fiber, copper, radio, satellite
  - transmission rate = *bandwidth*
- *Routers:* forward packets (chunks of data)



"Cool" Internet Appliances

IP picture frame
http://www.ceiva.com/

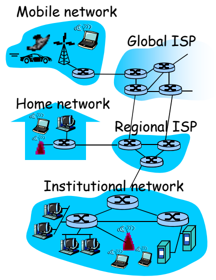
Web-enabled toaster + weather forecaster

World's smallest Web server
http://www-ccs.cs.umass.edu/~shri/iPic.html

Internet phones

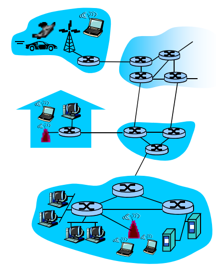
What's the Internet: "nuts and bolts" view

- *Protocols* control sending, receiving of msgs
  - e.g., TCP, IP, HTTP, Skype, Ethernet
- *Internet:* "network of networks"
  - loosely hierarchical
- Internet standards
  - Developed by IETF: Internet Engineering Task Force
  - IETF produces RFCs: Request for comments
    - More than 5000!
  - IEEE for links (e.g. 802.11)



What's the Internet: a Service View

- Communication *infrastructure* enables distributed applications:
  - Web, VoIP, email, games, e-commerce, file sharing
- Communication services provided to apps:
  - Reliable data delivery from source to destination
  - "best effort" (unreliable) data delivery



What's a Protocol?

Human protocols:

- "What's the time?"
  - Say "hi" first
- "I have a question"
  - Raise hand first

… specific msgs sent

… specific actions taken when msgs received, or other events

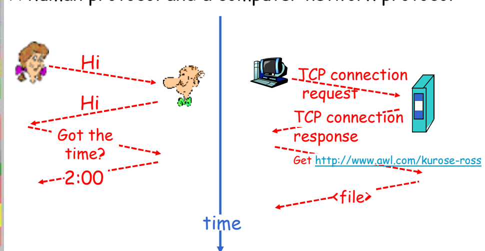
Network protocols:

- Machines rather than humans
- All communication activity in Internet governed by protocols

Protocols: 1) define format, 2) order of msgs sent and received among network entities, and 3) actions taken on msg transmission, receipt

What's a Protocol?

A human protocol and a computer network protocol:



Q: Other human protocols? Other network protocols?

Chapter 1: Roadmap

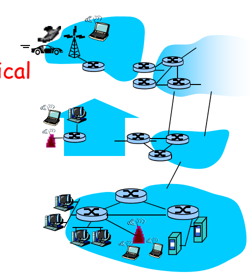
1.1 What *is* the Internet?
1.2 Network edge
- end systems, access networks, links

1.3 Network core
- circuit switching, packet switching, network structure

1.4 Delay, loss and throughput in packet-switched networks
1.5 Protocol layers, service models
1.6 Networks under attack: security
1.7 History

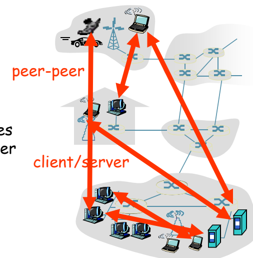
A Closer Look at Network Structure

- Network edge: applications and hosts
- Access networks, physical media: wired, wireless communication links
- Network core
  - interconnected routers
  - network of networks



The Network Edge

- End systems (hosts):
  - run application programs
  - e.g. Web, email
  - at "edge of network"
- Client/server model
  - client host requests, receives service from always-on server
  - e.g. Web browser/server; email client/server
- Peer-peer model:
  - minimal (or no) use of dedicated servers
  - e.g. Skype, BitTorrent



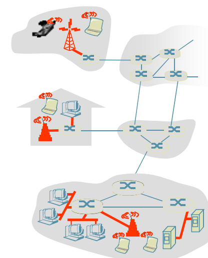
Access Networks and Physical Media

Q: How to connect end systems to edge router?

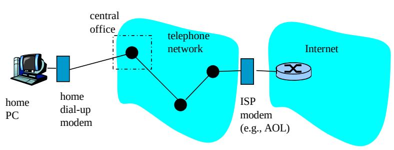
- residential access nets
- institutional access networks (school, company)
- mobile access networks

Keep in mind:

- bandwidth (bits per second) of access network?
- shared or dedicated?

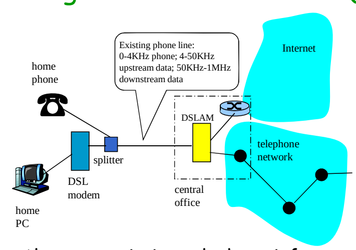


Dial-up Modem



- Uses existing telephony infrastructure
  - Home is connected to central office
- up to 56Kbps direct access to router (often less)
- Can't surf and phone at same time: not "always on"

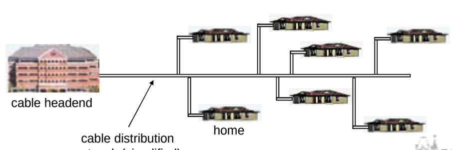
Digital Subscriber Line (DSL)



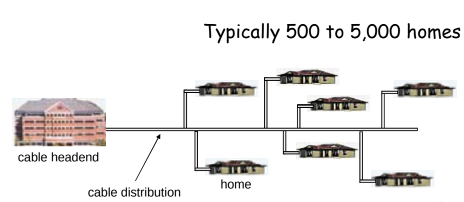
- Also uses existing telephone infrastruture
- up to 1 Mbps upstream (today typically < 256 kbps)
- up to 8 Mbps downstream (today typically < 1 Mbps)
- dedicated physical line to telephone central office

Cable Modems

- Does not use telephone infrastructure
  - Instead uses cable TV infrastructure
- HFC: hybrid fiber coax
  - asymmetric: up to 30 Mbps downstream, 2 Mbps upstream
- Network of cable and fiber attaches homes to ISP router
  - Homes share access to router (500 to 5,000 homes)
  - Unlike DSL, which has dedicated access



Cable Network Architecture: Overview



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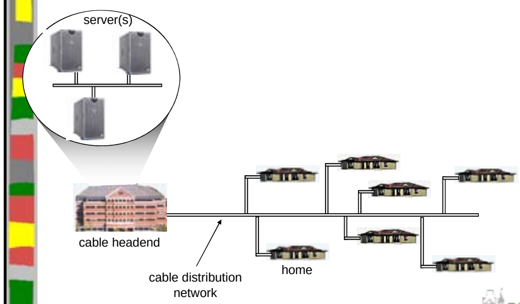

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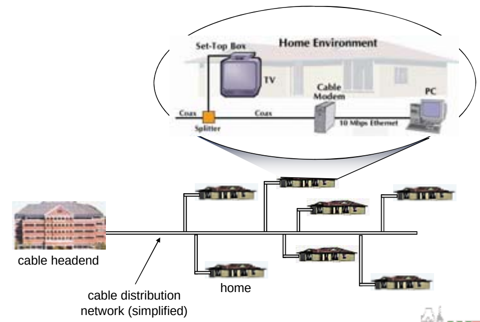
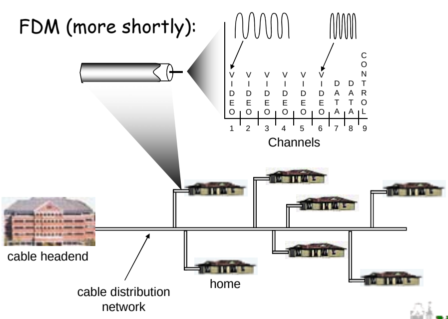
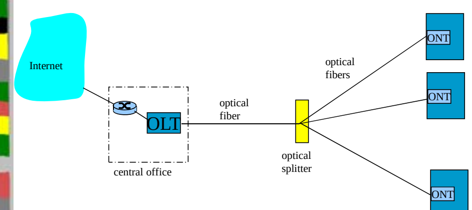


Cable Network Architecture: Overview

FDM (more shortly):

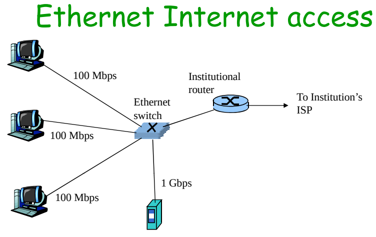


Fiber to the Home



- Optical links from central office to the home
- Much higher Internet rates; fiber also carries television and phone services

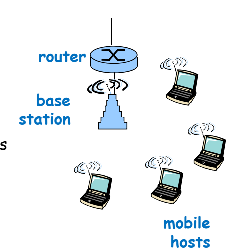
Ethernet Internet access



- Typically used in companies, universities, etc
- 10 Mbps, 100 Mbps, 1 Gbps, 10 Gbps Ethernet
- Today, end-systems typically connect into Ethernet switch

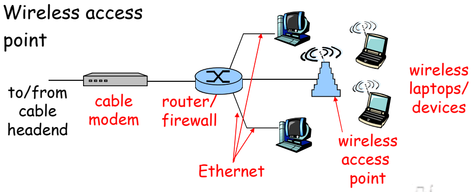
Wireless Access Networks

- *Shared* wireless access network connects end system to router
  - via base station, aka "access point" (AP)
- Wireless LANs:
  - 802.11b/g (WiFi): 11 or 54 Mbps
- Wider-area wireless access
  - Provided by telco operator
  - ~1Mbps over cell phone system (EVDO, HSDPA)
  - next up (?): WiMAX (10's Mbps) over wide area



Home Networks

Typical home network components:
- DSL or Cable modem
- Router/firewall/NAT
- Ethernet
- Wireless access point



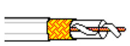
Physical Media: Twisted Pair

- Bit: propagates between transmitter/rcvr pairs
- Physical link: what lies between transmitter & receiver
- Guided media:
  - signals propagate in solid media: copper, fiber, coax
- Unguided media:
  - signals propagate freely, e.g., radio

Twisted Pair (TP)
- Two insulated copper wires
  - Category 3: traditional phone wires, 10 Mbps Ethernet
  - Category 5: 100 Mbps Ethernet

Physical Media: Coax, Fiber

Coaxial cable:
- Two concentric copper conductors
- Bidirectional
- Baseband:
  - single channel on cable
  - legacy Ethernet
- Broadband:
  - multiple channels on cable
  - HFC



Fiber optic cable:
- glass fiber carrying light pulses, each pulse a bit
- high-speed operation:
  - high-speed point-to-point transmission (e.g., 10's-100's Gps)
- low error rate: repeaters spaced far apart ; immune to electromagnetic noise

Physical Media: Radio

- Signal carried in electromagnetic spectrum
- No physical "wire"
- Bidirectional
- Propagation environment effects:
  - reflection
  - obstruction by objects
  - interference

Radio link types:
- terrestrial microwave
  - up to 45 Mbps channels
- LAN (e.g., Wifi)
  - 802.11b - 11 Mbps
  - 801.11g - 54 Mbps
- Wide-area (e.g., cellular)
  - 3G cellular ~ 1 Mbps
- Satellite
  - Kbps to 45Mbps channel (or multiple smaller channels)
  - 270 msec end-end delay
  - geosynchronous versus low altitude

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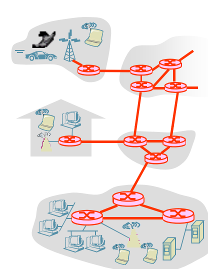
The Network Core

- Mesh of interconnected routers
- ***The*** fundamental question: how is data transferred through net?
  - circuit switching: dedicated circuit per call
    - E.g. telephone
  - packet-switching: data sent thru net in discrete "chunks"
    - E.g. postal mail

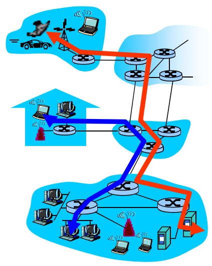


Q: human analogies of reserved resources (circuit switching) versus on-demand allocation (packet-switching)?

Network Core: Circuit Switching

End-end resources reserved for "call"

- link bandwidth, switch capacity
- dedicated resources: no sharing
- circuit-like (guaranteed) performance
- call setup required



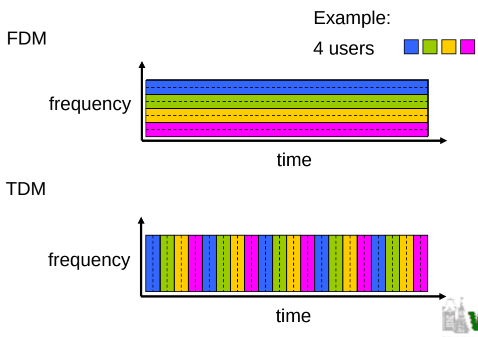
Network Core: Circuit Switching

network resources (e.g., bandwidth) divided into "pieces"

- pieces allocated to calls
- resource piece *idle* if not used by owning call *(no sharing)*

- dividing link bandwidth into "pieces"
  - frequency division multiplexing (FDM)
    - Radio frequency 88-108 MHz
    - Phone 4kHz
  - time division multiplexing (TDM)

Circuit Switching: FDM and TDM



Numerical Example

- How long does it take to send a file of 80 Kbytes from host A to host B over a circuit-switched network?
  - All links are 1.536 Mbps
  - Each link uses TDM with 24 slots/sec
  - 500 msec to establish end-to-end circuit

Let's work it out!

Numerical Example: Solution

- 80 Kbytes is 640,000 Kbits
  - NOTE: networks in bits, end systems in bytes
  - NOTE: 1 Kbyte = 1024 bytes, 1Kbit = 1000 bits
- Each circuit has a rate of 1.536 / 24 = 64 Kbps
- So, it takes 640,000 bits / 64 Kbps = 10 seconds to transmit the file
- Need to add the circuit establishment time (½ second)
- So, 10.5 seconds

Network Core: Packet Switching

Each end-end data stream divided into *packets*

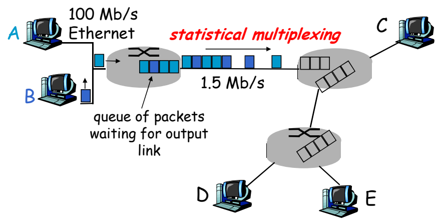
- user A, B packets *share* network resources
- each packet uses full link bandwidth
- resources used *as needed*

Bandwidth division into "pieces"
Dedicated allocation
Resource reservation

Resource contention:

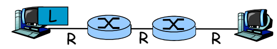
- aggregate resource demand can exceed amount available
- congestion: packets queue, wait for link use
- store and forward: packets move one hop at a time
  - Node receives complete packet before forwarding

Packet Switching: Statistical Multiplexing



Sequence of A & B packets does not have fixed pattern, bandwidth shared on demand ➔ ***statistical multiplexing***.

Packet-switching: Store-and-Forward



- Takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- *Store and forward:* entire packet must arrive at router before it can be transmitted on next link
- delay = 3 L/R (assuming zero propagation delay)

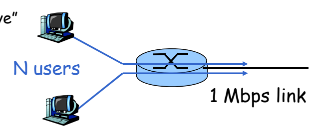
Example:
- 3 hops (end plus 2 routers)
- L = 7.5 Mbits
- R = 1.5 Mbps
- transmission delay
  - = 3 * 7.5 / 1.5
  - = 15 sec

more on delay shortly …

Packet Switching versus Circuit Switching

Packet switching can allow more users to use network!

- 1 Mb/s link
- Each user:
  - 100 kb/s when "active"
  - active 10% of time
- *circuit-switching:*
  - 10 users
- *packet switching:*
  - With 35 users, probability 10+ active at same time is less than .0004



Q: how did we get value 0.0004?

Packet Switching versus Circuit Switching

Packet switching can give individual users better performance!

- Consider
  - 3 users,
  - TDM with 1000 bit slot, 1 slot per 10 msec
- Users quiet, then one user 1000 1kbit packets
- With TDM, will take 10 seconds to transmit
- With packet switch, user can take all (1 Mbps) and transmit in about 1 second

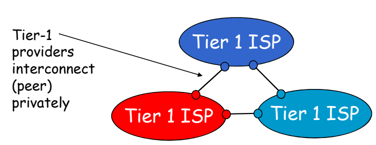
Packet Switching versus Circuit Switching

Is packet switching a "slam dunk" winner?

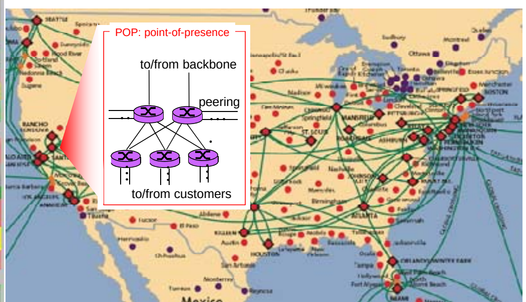
- Great for bursty data
  - Resource sharing
  - Simpler, no call setup
- But… can have excessive congestion: packet delay and loss
  - Protocols needed for reliable data transfer, congestion control
- Q: How to provide circuit-like behavior?
  - Bandwidth guarantees needed for audio/video apps
  - Still an unsolved problem (chapter 7)

Internet structure: Network of Networks

- Roughly hierarchical
- At center: "tier-1" ISPs (e.g., Verizon, Sprint, AT&T, Cable and Wireless), national/international coverage
  - treat each other as equals

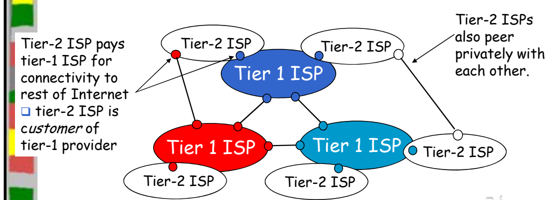


Tier-1 ISP: e.g., Sprint



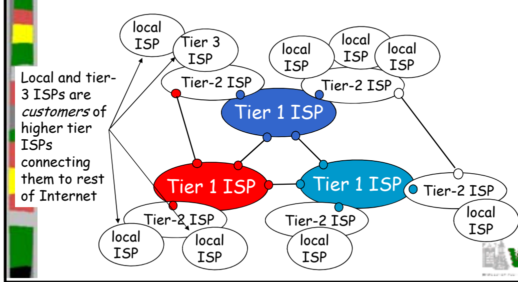
Internet structure: network of networks

- "Tier-2" ISPs: smaller (often regional) ISPs
  - Connect to one or more tier-1 ISPs, possibly other tier-2 ISPs



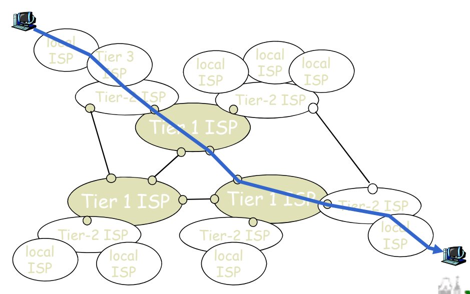
Internet structure: network of networks

- "Tier-3" ISPs and local ISPs
  - last hop ("access") network (closest to end systems)



Internet structure: network of networks

- a packet passes through many networks!



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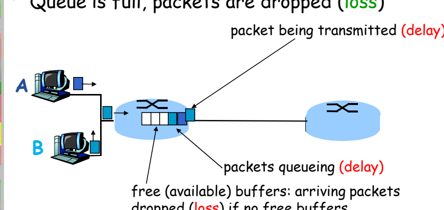
How do Loss and Delay Occur?

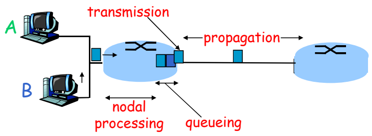
Packets *queue* in router buffers

- packet arrival rate to link exceeds output link capacity
- packets queue, wait for turn (delay)
- Queue is full, packets are dropped (loss)



Four Sources of Packet Delay

- 1. Nodal processing:
  - check bit errors
  - determine output link
- 2. Queueing
  - time waiting at output link for transmission
  - depends on congestion level of router

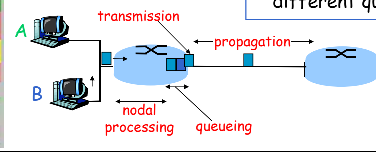


Delay in Packet-switched Networks

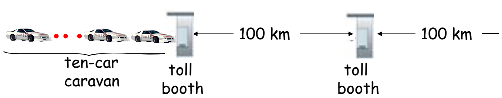
3. Transmission delay:
- R=link bandwidth (bps)
- L=packet length (bits)
- time to send bits into link = L/R

4. Propagation delay:
- d = length of physical link
- s = propagation speed in medium (~2×10^8 m/sec)
- propagation delay = d/s

Note: s and R are *very* different quantities!

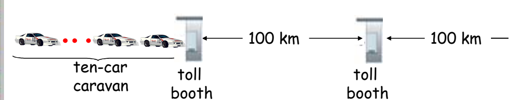


Caravan Analogy



- Car~bit; caravan~packet
- Cars "propagate" at 100 km/hr
- Toll booth takes 12 sec to service car (transmission time)
- Q: How long until caravan is lined up before 2nd toll booth?
- Time to "push" entire caravan through toll booth onto highway = 12*10 = 120 sec
- Time for last car to propagate from 1st to 2nd toll both: 100km/(100km/hr)= 1 hr
- A: 62 minutes

Caravan Analogy (more)



- Cars now "propagate" at 1000 km/hr
- Toll booth now takes 1 min to service a car
- Q: Will cars arrive to 2nd booth before all cars serviced at 1st booth?
- Yes! After 7 min, 1st car at 2nd booth and 3 cars still at 1st booth.
- 1st bit of packet can arrive at 2nd router before packet is fully transmitted at 1st router!
  - (See Ethernet Applet at book Web site)

Nodal delay

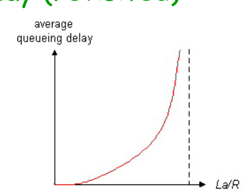
$$d_{nodal} = d_{proc} + d_{queue} + d_{trans} + d_{prop}$$

- d_{proc} = processing delay
  - typically a few microsecs or less
- d_{queue} = queuing delay
  - depends on congestion
- d_{trans} = transmission delay
  - = L/R, significant for low-speed links
- d_{prop} = propagation delay
  - a few microsecs to hundreds of msecs

Queueing Delay (revisited)

- R=link bandwidth (bps)
- L=packet length (bits)
- a=average packet arrival rate

traffic intensity = La/R



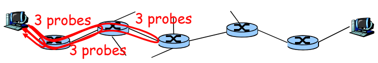
La/R ~ 0: average queueing delay small
La/R -> 1: delays become large
La/R > 1: more "work" arriving than can be serviced, average delay infinite!

“Real” Internet Delays and Routes

- What do “real” Internet delay & loss look like?

<u>Traceroute program:</u> provides delay measurement from source to router along end-end Internet path towards destination. For all *i*:

- sends three packets that will reach router *i* on path towards destination
- router *i* will return packets to sender
- sender times interval between transmission and reply.



“Real” Internet Delays and Routes

traceroute: gaia.cs.umass.edu to www.eurecom.fr

Three delay measurements from gaia.cs.umass.edu to cs-gw.cs.umass.edu

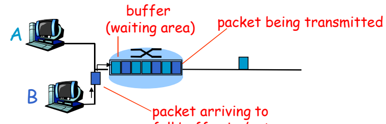
```
1  cs-gw (128.119.240.254)  1 ms  1 ms  2 ms
2  border1-rt-fa5-1-0.gw.umass.edu (128.119.3.145)  1 ms  1 ms  2 ms
3  cht-vbns.gw.umass.edu (128.119.3.130)  6 ms 5 ms 5 ms
4  jn1-at1-0-0-19.wor.vbns.net (204.147.132.129)  16 ms 11 ms 13 ms
5  jn1-so7-0-0-0.wae.vbns.net (204.147.136.136)  21 ms 18 ms 18 ms
6  abilene-vbns.abilene.ucaid.edu (198.32.11.9)  22 ms  18 ms  22 ms
7  nycm-wash.abilene.ucaid.edu (198.32.8.46)  22 ms  22 ms  22 ms
8  62.40.103.253 (62.40.103.253)  104 ms 109 ms 106 ms
9  de2-1.de1.de.geant.net (62.40.96.129)  109 ms 102 ms 104 ms
10  de.fr1.fr.geant.net (62.40.96.50)  113 ms 121 ms 114 ms
11  renater-gw.fr1.fr.geant.net (62.40.103.54)  112 ms  114 ms  112 ms
12  nio-n2.cssi.renater.fr (193.51.206.13)  111 ms  114 ms  116 ms
13  nice.cssi.renater.fr (195.220.98.102)  123 ms  125 ms  124 ms
14  r3t2-nice.cssi.renater.fr (195.220.98.110)  126 ms  126 ms  124 ms
15  eurecom-valbonne.r3t2.ft.net (193.48.50.54)  135 ms  128 ms  133 ms
16  194.214.211.25 (194.214.211.25)  126 ms  128 ms  126 ms
17  * * *
18  * * *
19  fantasia.eurecom.fr (193.55.113.142)  132 ms  128 ms  136 ms
```

trans-oceanic link (hops 7 → 8)

* means no response (probe lost, router not replying)

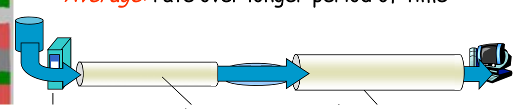
Packet Loss

- Queue (aka buffer) preceding link in buffer has finite capacity
- Packet arriving to full queue dropped (aka lost)
- Lost packet may be retransmitted by previous node, by source end system, or not at all



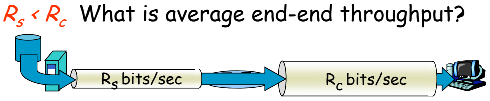
Throughput

- *Throughput:* rate (bits/time unit) at which bits transferred between sender/receiver
  - *Instantaneous:* rate at given point in time
  - *Average:* rate over longer period of time

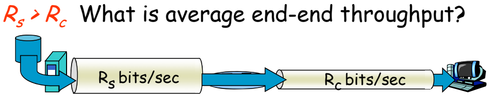


Throughput (more)

- $R_s < R_c$ What is average end-end throughput?



- $R_s > R_c$ What is average end-end throughput?

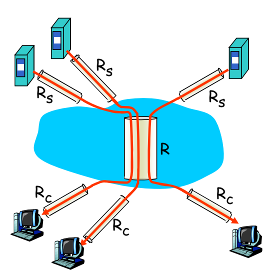


bottleneck link

link on end-end path that constrains end-end throughput

Throughput: Internet Scenario

- Per-connection end-end throughput: min(R_c,R_s,R/10)
- In practice: R_c or R_s is often bottleneck
  - “last mile” connection



10 connections (fairly) share backbone bottleneck link R bits/sec

Chapter 1: Roadmap

Protocol "Layers"

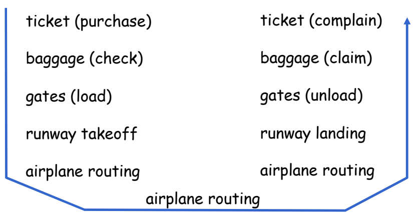
Networks are complex!

- Many "pieces":
  - hosts
  - routers
  - links of various media
  - applications
  - protocols
  - hardware, software

Question:
Is there any hope of *organizing* structure of network?

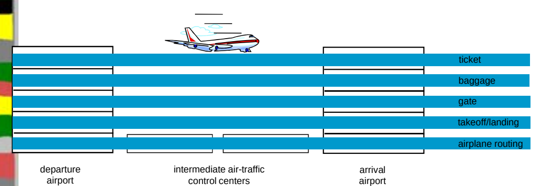
Or at least our discussion of networks?

Organization of air travel



- a series of steps

Layering of Airline Functionality



Layers: each layer implements a service
- via its own internal-layer actions
- relying on services provided by layer below

Why Layering?

Dealing with complex systems:
- explicit structure allows identification, relationship of complex system's pieces
  - layered reference model for discussion
- modularization eases maintenance, updating of system
  - change of implementation of layer's service transparent to rest of system
  - e.g., change in gate procedure doesn't affect rest of system
- layering considered harmful?
  - May need info from other layer (e.g. rate)
  - May be redundant functions (e.g. error check)

Internet Protocol Stack

- application: supporting network applications
  - FTP, SMTP, HTTP
- transport: process-process data transfer (segment)
  - TCP, UDP
- network: routing of datagrams from source to destination
  - IP, routing protocols
- link: data transfer between neighboring network elements (frames)
  - PPP, Ethernet
- physical: bits "on the wire"

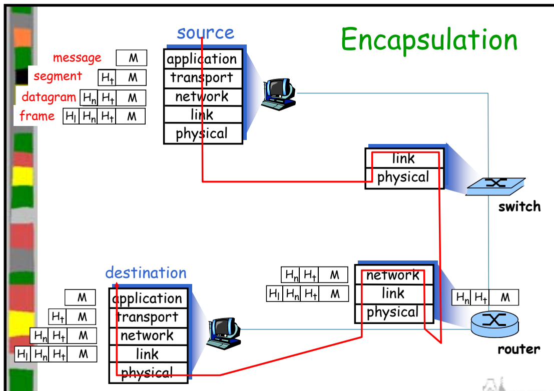
application
transport
network
link
physical

ISO/OSI Reference Model

- presentation: allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- *session:* synchronization, checkpointing, recovery of data exchange
- Internet stack "missing" these layers!
  - these services, *if needed,* must be implemented in application
  - needed?

application
presentation
session
transport
network
link
physical

Encapsulation



Chapter 1: Roadmap

1.1 What *is* the Internet?
1.2 Network edge
- end systems, access networks, links

1.3 Network core
- circuit switching, packet switching, network structure

1.4 Delay, loss and throughput in packet-switched networks
1.5 Protocol layers, service models
1.6 Networks under attack: security
1.7 History

Network Security

- The field of network security is about:
  - how bad guys can attack computer networks
  - how we can defend networks against attacks
  - how to design architectures that are immune to attacks
- Internet not originally designed with (much) security in mind
  - *original vision:* "a group of mutually trusting users attached to a transparent network" ☺
  - Internet protocol designers playing "catch-up"
  - Security considerations in all layers!

Bad guys can put malware into hosts via Internet

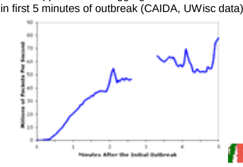
- Malware can get in host from a virus, worm, or trojan horse.
  - (More on next slide)
- Spyware malware can record keystrokes, Web sites visited, upload info to collection site
- Infected host can be enrolled in a botnet, used for spam and DDoS attacks
- Malware is often self-replicating: from an infected host, seeks entry into other hosts

Bad guys can put malware into hosts via Internet

- Trojan horse
  - Hidden part of some otherwise useful software
  - Today often on a Web page (Active-X, plugin)
- Virus
  - infection by receiving object (e.g., e-mail attachment), actively executing
  - self-replicating: propagate itself to other hosts, users

- Worm:
  - infection by passively receiving object that gets itself executed
  - self- replicating: propagates to other hosts, users

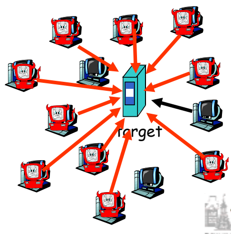
Sapphire Worm: aggregate scans/sec in first 5 minutes of outbreak (CAIDA, UWisc data)



Bad guys can attack servers and network infrastructure

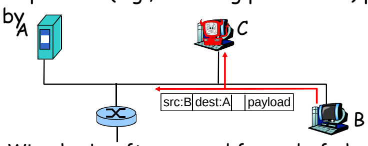
- Denial of service (DoS): attackers make resources (server, bandwidth) unavailable to legitimate traffic by overwhelming resource with bogus traffic

1. select target
2. break into hosts around the network (see botnet)
3. send packets toward target from compromised hosts



The bad guys can sniff packets

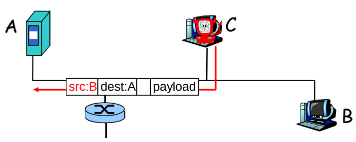
Packet sniffing:
- broadcast media (shared Ethernet, wireless)
- promiscuous network interface reads/records all packets (e.g., including passwords!) passing by



- Wireshark software used for end-of-chapter labs is a (free) packet-sniffer

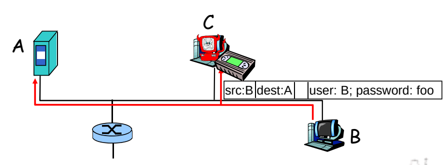
The bad guys can use false source addresses

- *IP spoofing:* send packet with false source address



The bad guys can record and playback

- *record-and-playback*: sniff sensitive info (e.g., password), and use later
  - password holder *is* that user from system point of view



Network Security

- A bit in this course
- Chapter 8: focus on security

Chapter 1: roadmap

Internet History

1961-1972: Early packet-switching principles

- 1961: Kleinrock - queueing theory shows effectiveness of packet-switching
- 1964: Baran - packet-switching in military nets
- 1967: ARPAnet conceived by Advanced Research Projects Agency
- 1969: first ARPAnet node operational
- 1972:
  - ARPAnet public demonstration
  - NCP (Network Control Protocol) first host-host protocol
  - first e-mail program
  - ARPAnet has 15 nodes

Internet History

1972-1980: Internetworking, new and proprietary nets

- 1970: ALOHAnet satellite network in Hawaii
- 1974: Cerf and Kahn - architecture for interconnecting networks
- 1976: Ethernet at Xerox PARC
- ate70's: proprietary architectures: DECnet, SNA, XNA
- late 70's: switching fixed length packets (ATM precursor)
- 1979: ARPAnet has 200 nodes

Cerf and Kahn's internetworking principles:
- minimalism, autonomy - no internal changes required to interconnect networks
- best effort service model
- stateless routers
- decentralized control

define today's Internet architecture

Internet History

1980-1990: new protocols, a proliferation of networks

- 1983: deployment of TCP/IP
- 1982: smtp e-mail protocol defined
- 1983: DNS defined for name-to-IP-address translation
- 1985: ftp protocol defined
- 1988: TCP congestion control
- new national networks: Csnet, BITnet, NSFnet, Minitel
- 100,000 hosts connected to confederation of networks

Internet History

1990, 2000's: commercialization, the Web, new apps

- Early 1990's: ARPAnet decommissioned
- 1991: NSF lifts restrictions on commercial use of NSFnet (decommissioned, 1995)
- early 1990s: Web
  - hypertext [Bush 1945, Nelson 1960's]
  - HTML, HTTP: Berners-Lee
  - 1994: Mosaic, later Netscape
  - late 1990's: commercialization of the Web

Late 1990's – 2000's:
- more killer apps: instant messaging, P2P file sharing
- network security to forefront
- est. 50 million host, 100 million+ users
- backbone links running at Gbps

Internet History

2007:
- ~500 million hosts
- Voice, Video over IP
- P2P applications: BitTorrent (file sharing) Skype (VoIP), PPLive (video)
- more applications: YouTube, gaming
- wireless, mobility

Introduction: Summary

Covered a "ton" of material!
- Internet overview
- what's a protocol?
- network edge, core, access network
  - packet-switching versus circuit-switching
  - Internet structure
- performance: loss, delay, throughput
- layering, service models
- security
- history

You now have:
- context, overview, "feel" of networking
- more depth, detail *to follow!*

Introduction

CS 3516 – Computer Networks