

**ABE: Providing a Low Delay
within Best Effort**

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Introduction

- Multimedia applications can perform well over a wide-range of loss (with repair)
- Delay often the major impediment for interactive multimedia applications
- Internet is “best-effort” with one QoS of traffic for all
 - DiffServ requires monitoring of classes
- Want to keep it simple, but add support for delay sensitive multimedia traffic
 - *Alternative Best Effort (ABE)*

Outline

- Introduction (done)
- The ABE Service (next)
- Implementation
- Simulation Results
- Related Work
- Conclusions

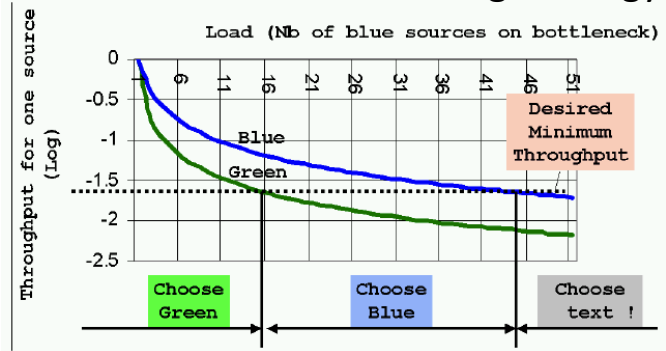
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- Introduction (done)
- The ABE Service (next)
 - Definition
 - Green does not hurt blue
 - Router requirements
 - Inter-working and Migration
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Definition

- ABE packets are either **green** or **blue**
 - (Neutral colors, **green** for “go”)
 - Application chooses to make packets **green**
 - Default is **blue** (most applications)
- **Green** packets get low, bounded delay
- **Green** does not hurt **blue**
 - **Blue** has same or better throughput even in presence of **green** traffic
- All ABE packets in same best-effort class
 - Traditional congestion control
 - All **blue** get more throughput than all **green**

Possible Packet Coloring Strategy



Assume: $utility(rate, delay) = 0$ if rate < min
 $utility(rate, delay) = \text{linear with delay}$ if rate > min

Discussion

- Interactive applications send mix of **blue** and **green**
 - “Probe” packets to determine region
- Traditional applications send all **blue**
 - Care more about throughput
- Note, says nothing about *TCP-friendly*
 - Still same problem as with best-effort
 - **Green** makes it no worse since doesn’t hurt **blue**
- Backbones have low delay, so ABE likely in peripheral/edge routers
- Delay bound offered depends upon hops
 - Assume 2-6 low-speed hops
 - Delay 100-150 msec total, maybe 50 for network
 - Per-hop delay about 5-20 msec

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Green Does Not Hurt Blue

- When there is **green** traffic in addition to traditional **blue** traffic, must have
 - Local transparency to **blue**
 - Throughput transparency to **blue**

Local Transparency to Blue

- Consider a traditional router that treats all packets equal (no ABE)
 - ABE router should have same delay as traditional router
 - If **blue** not dropped with traditional router, then not dropped with ABE router
- If TCP friendly:

$$\theta = \frac{s}{R\sqrt{\frac{2p}{3}} + 3t_1\sqrt{\frac{3p}{8}}p(1 + 32p^2)}$$

[11, Kurose SIGCOMM '98] – well-used model

- If reduce R, what might happen to throughput for **green**?
 - Need throughput transparency

Throughput Transparency to Blue

- If **green** flow is TCP friendly, should get less or equal throughput as **blue** flows
- Hard to implement exactly since hard to measure
 - Hard to measure TCP friendly, even!
 - Consider it to be a loose requirement
- Implement by making sure **green** has higher loss ratio

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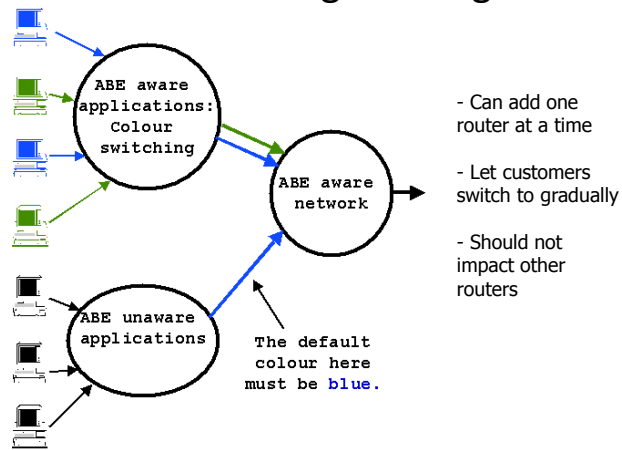
Router Requirements

- Provide low, bounded delay to **green**
- Provide local transparency to **blue**
- Provide throughput transparency to **blue**
- Preserve packet sequence within **blue** and **green**
 - Note, may be out of order across colors
- Keep **green** packet loss as low as possible
 - Make **green** attractive as possible

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Interworking and Migration



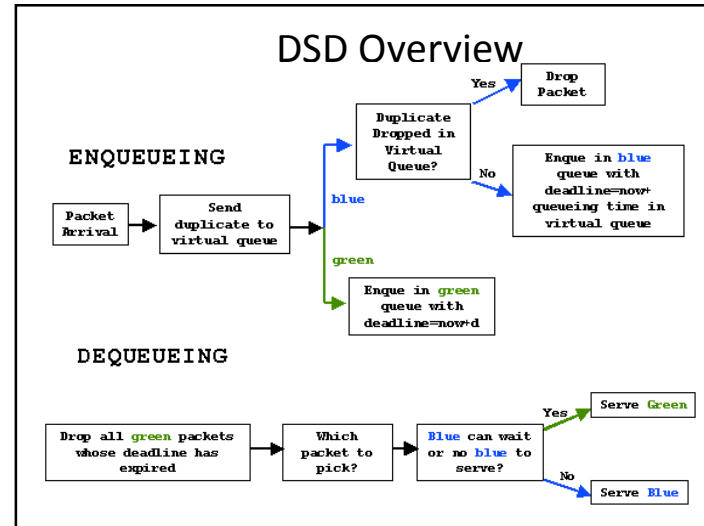
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- Introduction **(done)**
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- Implementation **(next)**
 - Duplicate Scheduling with Deadlines
 - Properties of (DSD)
- Simulation Results
- Related Work
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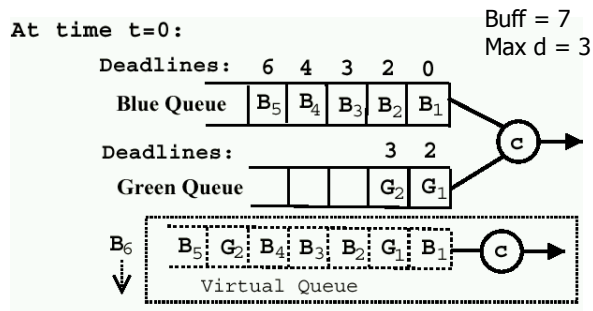
Implementation

- Could try modified FCFS:
 - For **blue**, enqueue normally
 - For **green**, drop if delay > max
 - (What is problem with only doing this?)
 - Instead, use separate queues
 - But still *work conserving*
 - Deadlines associated with each packet
 - Dequeue color that has earlier deadline
 - If both, use control function for fairness
- *Duplicate Scheduling with Deadlines (DSD)*

DSD Overview

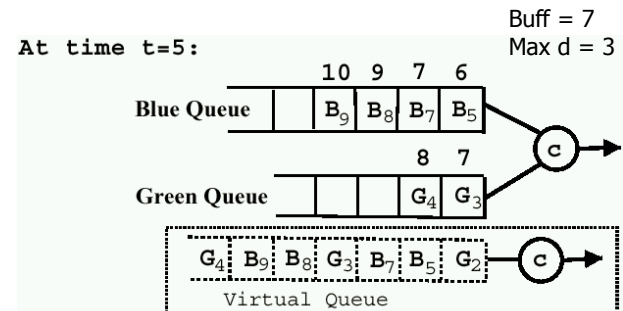


DSD Example (1 of 2)



Serve: B₁, G₁, B₂, B₃, B₄
 Drop: G₂ (deadline missed), B₆ (buffer full)

DSD Example (2 of 2)



Serve: G₃, B₅, B₇, G₄, B₈ and B₉

DSD Modifications

- Only enqueue green packet if length of green queue + blue packets with deadline less than d
 - (would not have enqueued G2, in example)
- If either can be served, if $[0,1] < g$ then pick green else blue
 - $g=1$ favor green, $g=0$ favor blue
 - ($g=1$ in example)
- Can also use active queue management (AQM, say RED) for congestion monitoring

Properties of DSD

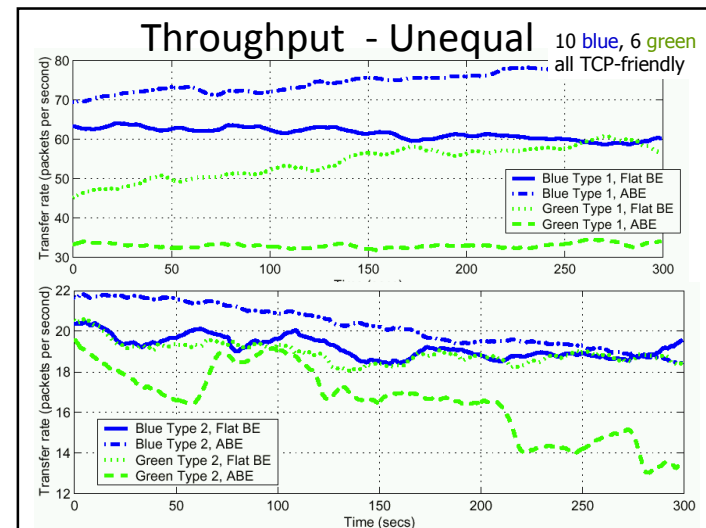
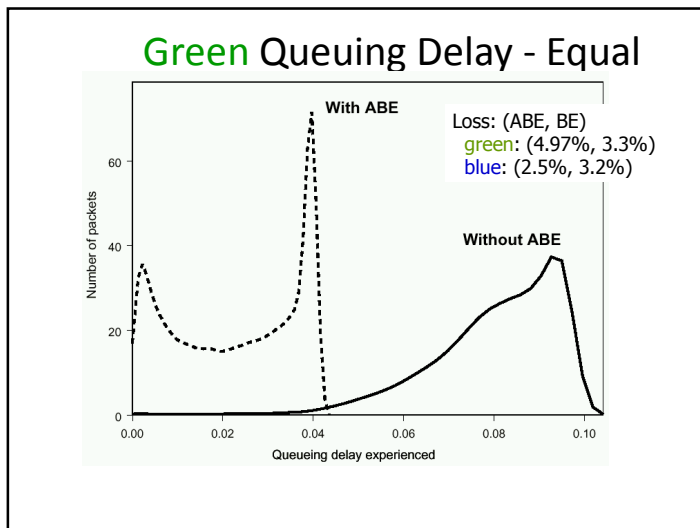
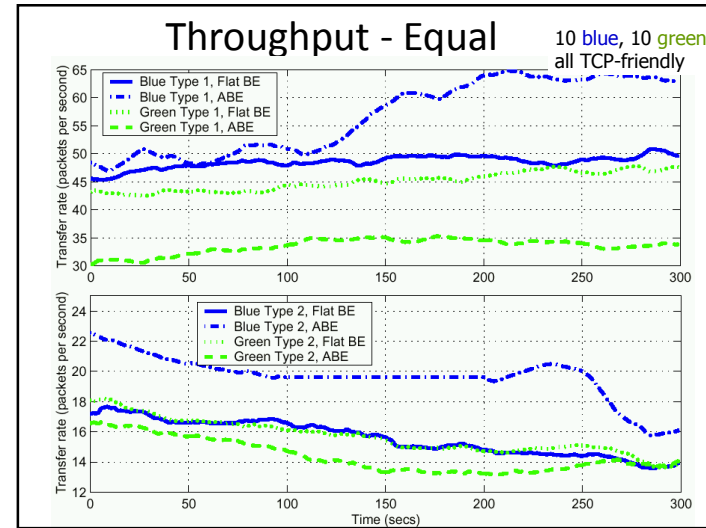
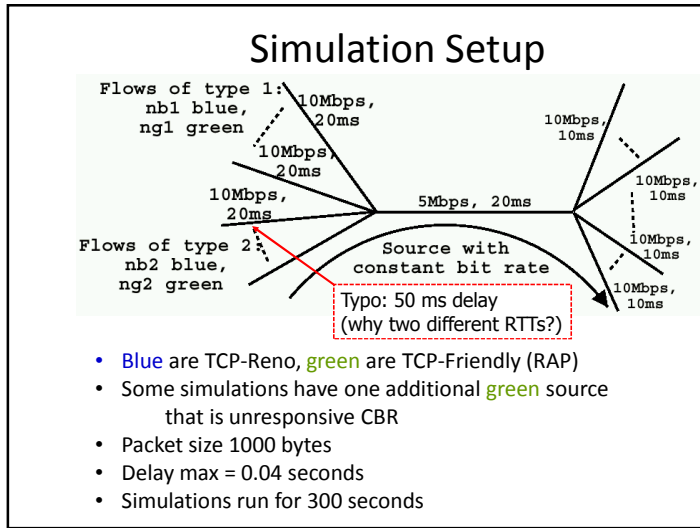
- Buffer always less than $Buff$ because of virtual queue
- All blue packets served by deadlines, so same as or earlier than best-effort
- All green packets served before d , else dropped

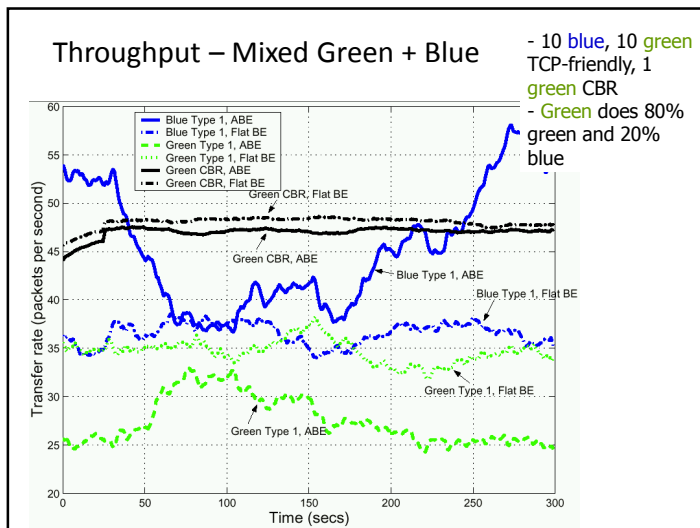
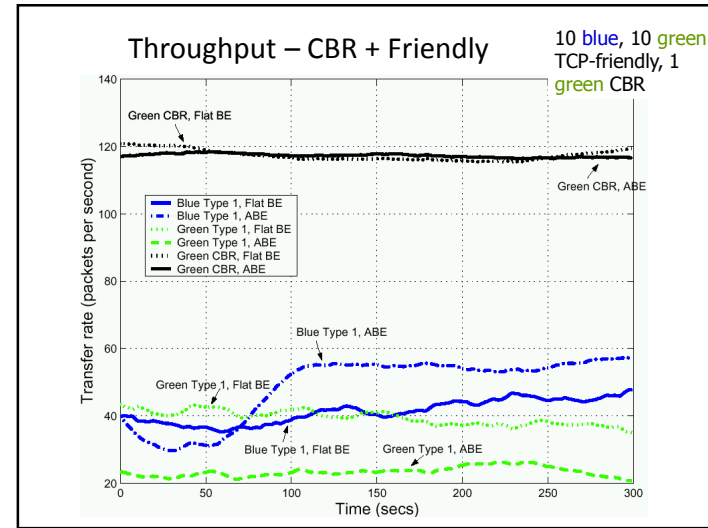
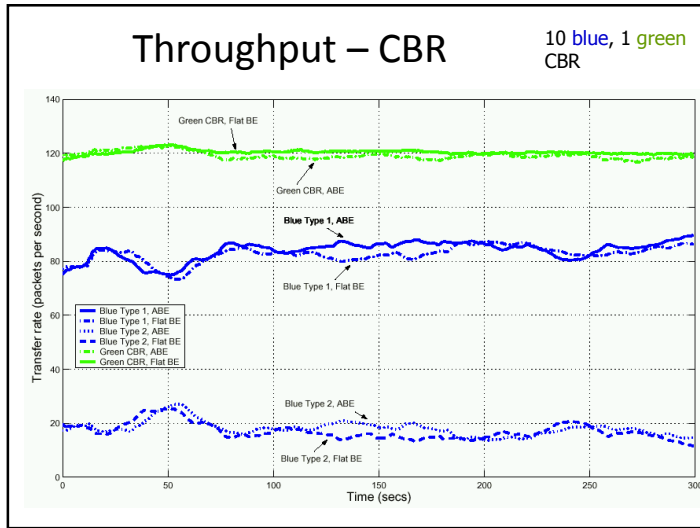
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Simulation

- Done in NS-2
- Show green does not hurt blue
- Show green benefits from low delay
- Show loss rates for both types
- Compare to reference condition, flat best-effort FCFS (droptail) router





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Related Work

- IntServ [21, 22]
 - Admission control plus reservation
 - Per-flow accounting and charging
 - Doesn't scale
 - May perform on edge only
- DiffServ
 - Aggregates (classes) of flows
 - Scales better
 - (Next chunk of related work)

Related Work

- DiffServ Low delay service
 - Crowcroft et al (also gets more throughput) [23]
 - EF provides low delay and low loss [25]
 - SIMA has level for how 'real-time' traffic is [26]
- DiffServ Low delay class
 - Dovrolis et al [27]
 - AF – Assured Forwarding [30]
- All require changes to existing price structures. Incremental deployment difficult.

Conclusion

- ABE
 - Supports low delay
 - No reservation or signaling required
- Choice of green or blue up to application
- One ABE implementation presented (DSD)
- Simulation and implementation suggest:
 - Green benefits from lower delay
 - Blue not harmed
 - Over a range of network and traffic conditions

Future Work?

Future Work

- Applications that use **green**
 - Adaptively
- PQ benefits of ABE to multimedia apps
- Implementation overhead of ABE
- More colors for more MM applications:
 - **dark green**, **light green**, **neon green** ...
- More colors for more **blue** applications
 - Web, Email, Telnet, File Transfer