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SEARCH TREE

Node: State in state tree

Root node: Top of state tree

Children: Nodes that can be reached from a given

node in 1 step (1 operator)

Expanding: Generating the children of a node

Open: Node not yet expanded

Closed: Node after expansion

Queue: Ordered list of open nodes

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SEARCH

BLIND SEARCH: Systematic Search

Depth-1st: Continue along current path looking

for goal

Breadth-1st: Expand all nodes at current level

before progressing to next level

Depth-limited Search: Depth-1st + depth-limit

Iterative Deepening Search: limit=0,limit=1,...

USING COST: g(n) = cost from start to n

Branch-and-bound (= Uniform Cost Search):

Select node n with best g(n).

USING HEURISTIC: h(n)=Estimate cost to a goal

Greedy Search: Select node n with best h(n)

A*: Select node *n* with best f(n) = g(n) + h(n)

IDA*: $A^* + f$ -cost limit.

Hill-Climbing: Depth-1st exploring best h(n) first

Simulated Annealing: Hill-Climbing + RandomWalk

Beam Search: Breadth-1st keeping only m nodes

with best h(n)'s per level

Al Lecture. Prof. Carolina Ruiz. Worcester Polytechnic Institute DEPTH—1st SEARCH DEPTH—1st (cont.) When to use If queue is empty then fail If head of queue is goal then succeed Else remove head of queue, expand it, place children in front of queue Prof. Carolina Ruiz. Worcester Polytechnic Institute DEPTH—1st (cont.) When to use All solutions at same depth All solutions at same depth Any solution will do Possibly fast	Page 2
DEPTH-1st SEARCH 1. Put start state onto queue 2. If queue is empty then fail 3. If head of queue is goal then succeed 4. Else remove head of queue, expand it, place children in front of queue 5. Recurse to 2 DEPTH-1st (cont.) When to use • Depth limited or known beforehand • All solutions at same depth • Any solution will do • Possibly fast When to avoid • Large or infinite subtrees	age 4
 2. If queue is empty then fail 3. If head of queue is goal then succeed 4. Else remove head of queue, expand it, place children in front of queue 5. Recurse to 2 Depth limited or known beforehand All solutions at same depth Any solution will do Possibly fast When to avoid Large or infinite subtrees 	
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. Else remove head of queue, expand it, place children in front of queue • Possibly fast When to avoid • Large or infinite subtrees	
in front of queue Recurse to 2 When to avoid Large or infinite subtrees	
• Large or infinite subtrees	
• Prefer shallow solution	
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AI Lecture. Prof. Carolina Ruiz. Worcester Polytechnic Instir BREADTH-1st SEARCH	tute	AI Lecture. Prof. Carolina Ruiz. Worcester Polytechnic Institute BREADTH-1st (Cont.)	
1. Put start state onto queue		When to use	
2. If queue is empty then fail		• Large or infinite search tree	
3. If head of queue is goal then succeed		• Solution depth unknown	
4. Else remove head of queue, expand it, place	e children	• Prefer shallow solution	
at end of queue		When to avoid	
5. Recurse to 2		• Very wide trees	
		• Generally slow	
		• May need a lot of space	

MODIFICATIONS TO DEPTH/BREADTH 1ST

Depth-limited Search:

Limit the total depth of the depth 1st search.

Iterative Deepening Search:

Repeat depth—limited search with limit 0, 1, 2, 3, ... until a solution is found.

Bidirectional Search:

Simultaneously search forward from initial state and backward from goal state until both paths meet.

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BRANCH-AND-BOUND (= UNIFORM-COST SEARCH)

- 1. Put start state onto queue
- 2. If queue is empty then fail
- 3. If head of queue is goal then succeed
- 4. Else
 - remove head of queue,
 - expand it,
 - place in queue, and
 - sort entire queue with least cost-so-far nodes in front
- 5. Recurse to 2

BRANCH-AND-BOUND SUMMARY

Advantages

- Optimal (when costs are non-negative)
- Complete

Disadvantages

• Can be inefficient

When to use

- Desire best solution
- Keep track of cost so far

When to avoid

- May not work with negative costs
- May be overly conservative
- Any solution will do

Potential improvement

• Dynamic Programming

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BRANCH-AND-BOUND + DYNAMIC PROG.

- 1. Put start state onto queue
- 2. If queue is empty then fail
- 3. If head of queue is goal then succeed
- 4. Else
 - remove head of queue,
 - expand it,
 - place in queue,
 - ★ remove redundant paths:Paths that reach the same node as other paths but are more expensive, and
 - sort entire queue with least cost-so-far nodes in front
- 5. Recurse to 2

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GREEDY SEARCH (= Winston's BEST-1st SEARCH)

- 1. Put start state onto queue
- 2. If queue is empty then fail
- 3. If head of queue is goal then succeed
- 4. Else
 - remove head of queue,
 - expand it,
 - place in queue, and
 - sort entire queue with least estimated-costto-goal nodes in front
- 5. Recurse to 2

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GREEDY SEARCH SUMMARY

Advantages

- Can be very efficient
- Paths found are likely to be short

Disadvantages

• Neither optimal nor complete

When to use

• Desire "short" solution

When to avoid

• When an optimal solution is required

\mathbf{A}^*

- 1. Put start state onto queue
- 2. If queue is empty then fail
- 3. If head of queue is goal then succeed
- 4. Else remove head of queue, expand it, place in queue, and **sort entire queue** with **least cost-so-far** + **estimated-cost-remaining** nodes in front
- 5. If multiple paths reach a common goal, keep only lowest cost-so-far path
- 6. Recurse to 2
- f(node) = g(node) + h(node), where
 - -f(node) = estimated total cost
 - -g(node) = cost-so-far to node
 - -h(node) = estimated-cost-remaining (heuristic).
- Properties of h:
 - Lower bound (\leq actual cost)
 - Nonnegative

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A* SUMMARY

Advantages

- Complete
- \bullet Optimal, when h is an underestimate
- Optimally efficient among all optimal search algorithms

Disadvantages

• Very high space complexity

When to use

- Desire best solution
- Keep track of cost so far
- Heuristic information available

When to avoid

• No good heuristics available

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HILL CLIMBING SEARCH		HILL CLIMBING SUMMARY		
1. Put start state onto queue		Advantages		
2. If queue is empty then fail		• Complete if backtracking is allowed (like in	Win-	
3. If head of queue is goal then succeed		ston's book) and the graph is finite		
4. Else remove head of queue, expand it, place childr	en	Disadvantages		
sorted by $h(n)$ in front of queue		• Not optimal		
5. Recurse to 2		When to use		
		• Depth limited or known beforehand		

 \bullet All solutions at same depth

 \bullet If optimal solution is required

 \bullet Large or infinite subtrees

• No good estimate • Difficult terrain

 \bullet Reliable estimate of remaining distance to goal

 \bullet Desire good solution

 \bullet Fast if good estimate

When to avoid

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Search Methods Page 17 Search Methods AI Lecture. Prof. Carolina Ruiz. Worcester Polytechnic Institute BEAM SEARCH BEAM SEARCH SUMMARY 1. Put start state onto queue Advantages 2. If queue is empty then fail • Saves space 3. If head of queue is goal then succeed Disadvantages 4. Else remove head of queue, expand it, place children • Neither optimal nor complete at end of queue When to use 5. If finishing a level, keep only w best nodes in queue • Large or infinite search tree 6. Recurse to 2 • Solution depth unknown • Prefer shallow solution • Possibly fast

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• No more than wb nodes stored

When to avoid

- Can't tell which solutions to prune
- Prefer conservative

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SEARCH STRATEGIES -

Completeness; Optimality; and Time and Space Complexity

Search	Complete?	Optimal?	Time	Space
Depth-1st	N	N	b^d	bd
Breadth-1st	Y	Y*	b^s	b^s
Depth-limited	N	N	b^l	bl
Iter. deepening	Y	Y*	b^s	bs
Branch-&-bound	Y	Y	b^s	b^s
Greedy	N	N	b^d	b^d
A*	Y	Y	exp	exp
Hill-climbing	N	N	dep	dep
Beam	N	N	ms	2m

(adapted from Russell & Norvig's book)

- Y*: Yes, IF cost of a path is equal to its length. Otherwise No.
- \bullet b: branching factor
- \bullet s: depth of the solution
- \bullet d: maximum depth of the search tree
- \bullet l: depth limit
- m: beam size
- ullet exp: exponential depending on heuristic h
- \bullet dep: depends on heuristic h

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Search Methods

SEARCH STRATEGIES IN WINSTON'S BOOK Summary

Depth 1st: Continue along current path looking for goal

Breadth 1st: Expand all nodes at current level before progressing to next level

Hill Climbing: Like depth 1st, but explore most promising children first

Beam: Like breadth 1st, but prune unpromising children

Best 1st: Expand best open node regardless of its depth

Branch-and-bound: Expand the least-cost-so-far node until goal reached

A*: Like branch-and-bound, but with heuristic information