

CS 525M – Mobile and Ubiquitous Computing Seminar

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Broadcast Disk Implementation

- Contents
 - Introduction
 - Design
 - Results
 - Conclusion
 - Future Work

Broadcast Disk Implementation

- Introduction
 - Implemented Multicast Broadcast disk
 - Client / Server
 - Cache Policy
 - Broadcast Disk Algorithm
 - Test Script

Broadcast Disk Implementation

- Design

- Implementation tools

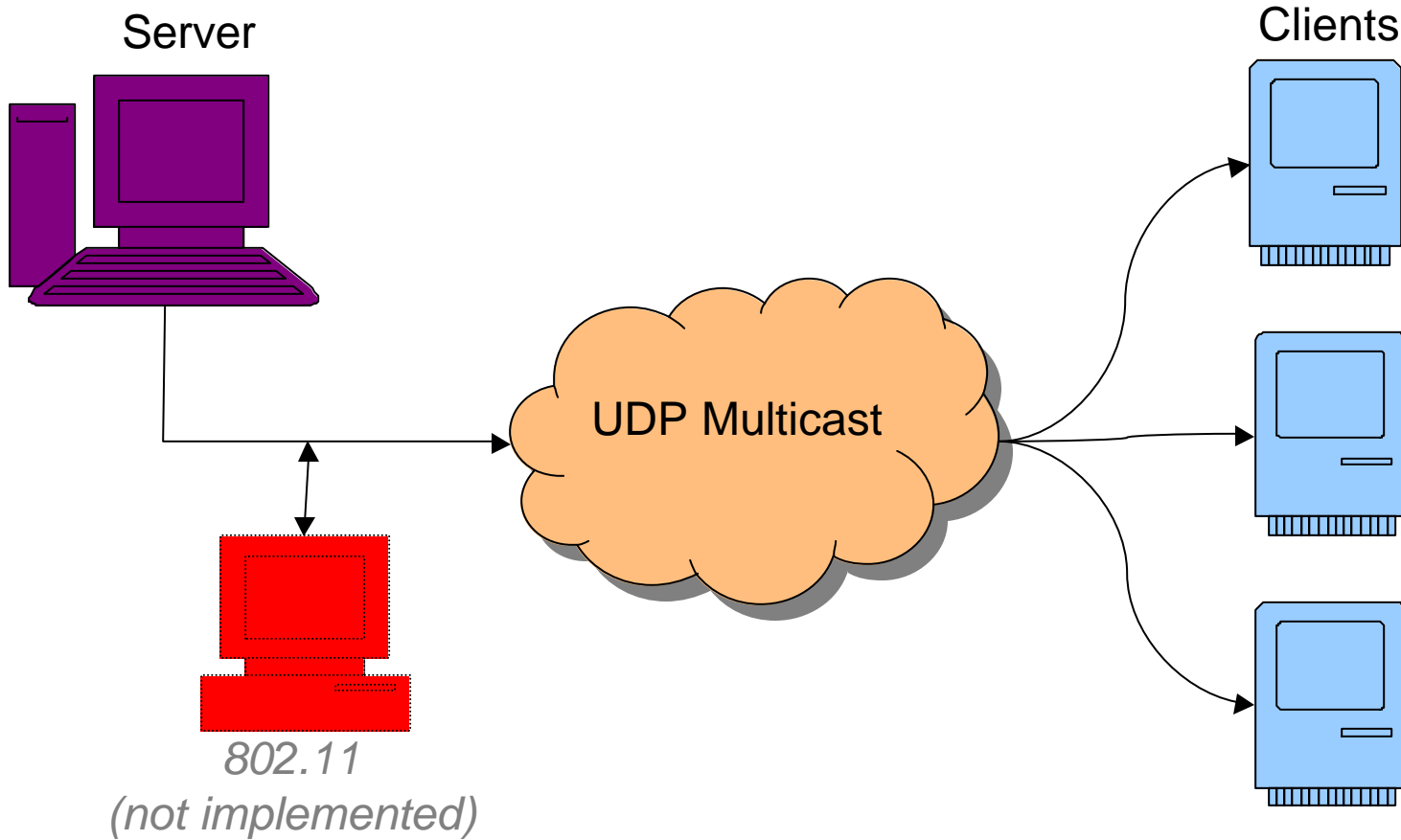
- UDP Packets

- Java Networking

- Eclipse IDE for development

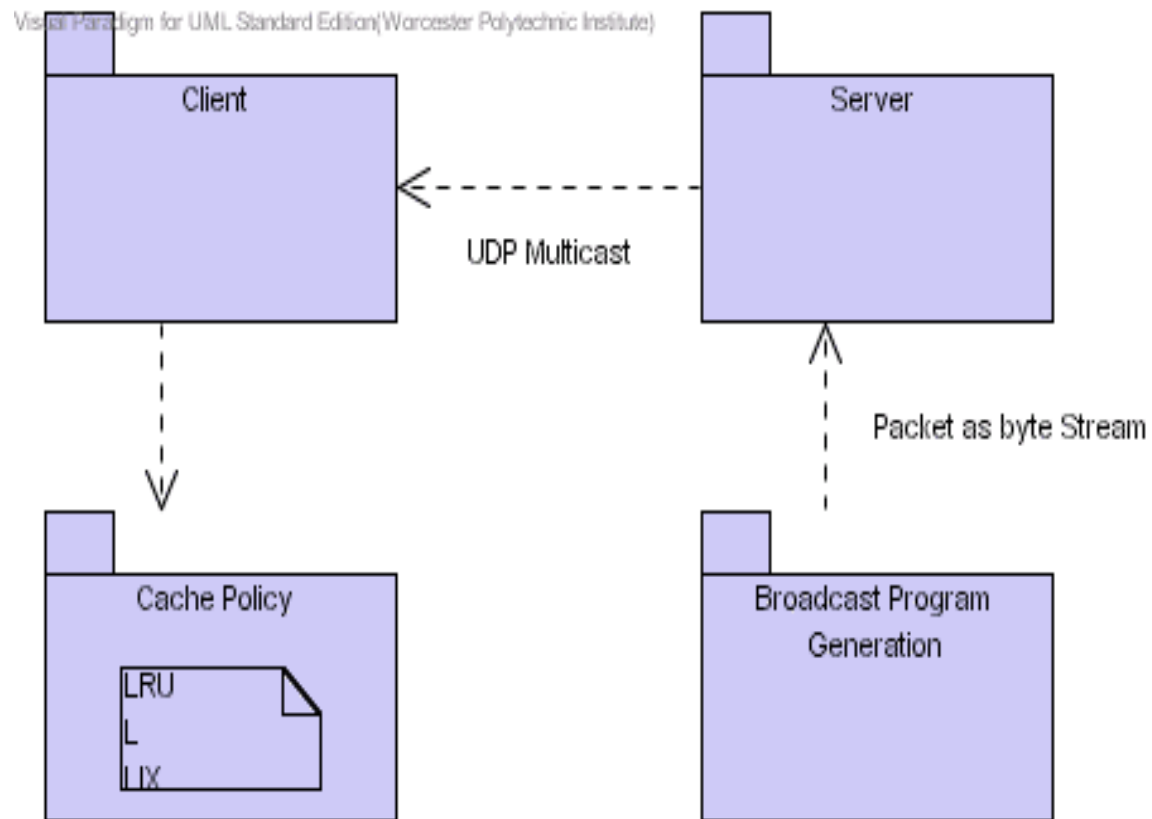
Broadcast Disk Implementation

Network Layout



Broadcast Disk Implementation

Design Diagram



Broadcast Disk Implementation

Design- Cache Policy

– Implemented

- LRU - linked list accessed pages are moved to the top of the list
- LIX - Link list for each broadcast disk smallest *lix* value of bottom pages ejected.
- L - like LIX except same frequency value for each disk

– Not Implemented

- P – highest access probability in cache
- PIX – lowest ratio of access probability to broadcast frequency

Broadcast Disk Implementation

Design- Cache Policy

– Lix Example

- $p_i = ? / (\text{currentTime} - t_i) + (1 - ?)p_i$
- $\text{lix} = p_i / \text{frequency of the page}$

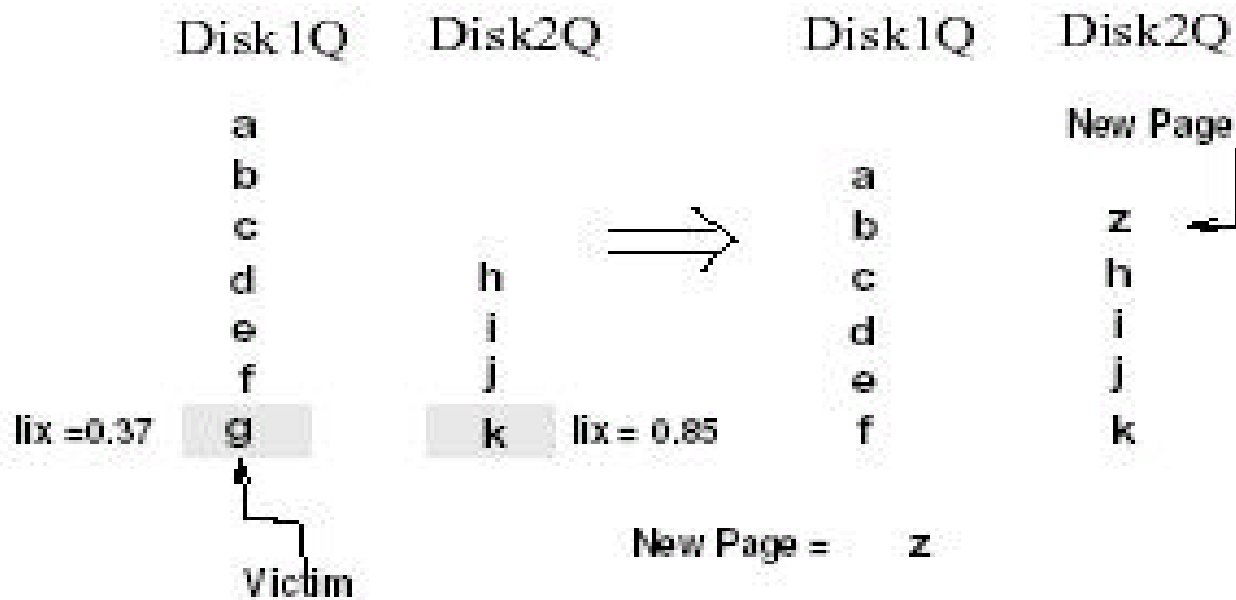


Figure 12: Page replacement in \hat{LIX}

Broadcast Disk Implementation

Design – Broadcast Program Generation

- Create page vector
- Apply noise values
- Calculates Chunks per disk
- Generate relative frequency per disk
- Interleave chunks while applying offset
- Generate page list

Broadcast Disk Implementation

Design – Client

- Executes as a java thread
- Listens on UDP multicast
- Uses three different caches or no cache
- Receives the following parameters:

<i>CacheSize</i>	Client cache size (in pages)
<i>ThinkTime</i>	Time between client page accesses (in broadcast units)
<i>AccessRange</i>	# of pages in range accessed by client
θ	Zipf distribution parameter
<i>RegionSize</i>	# of pages per region for Zipf distribution

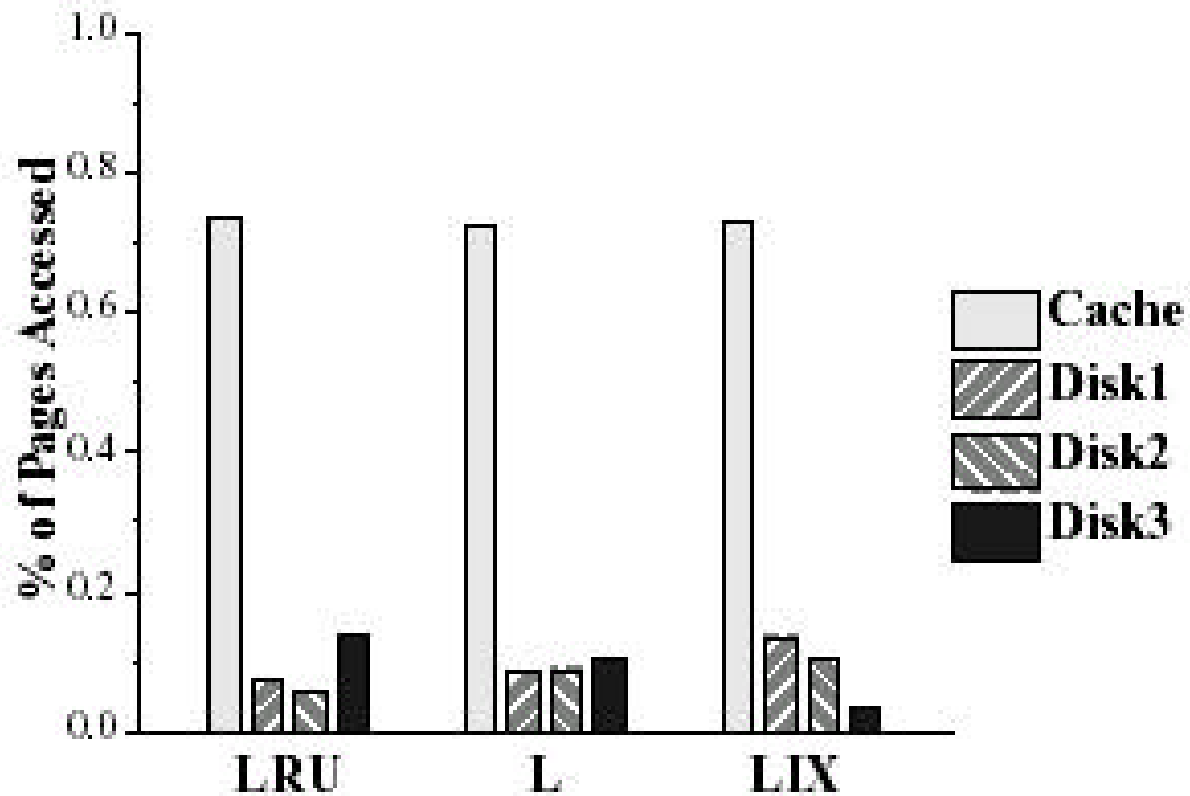
Broadcast Disk Implementation

Design – Server

- Executed as a thread.
- Sends UDP packets over multicast
- Receives the following parameters:

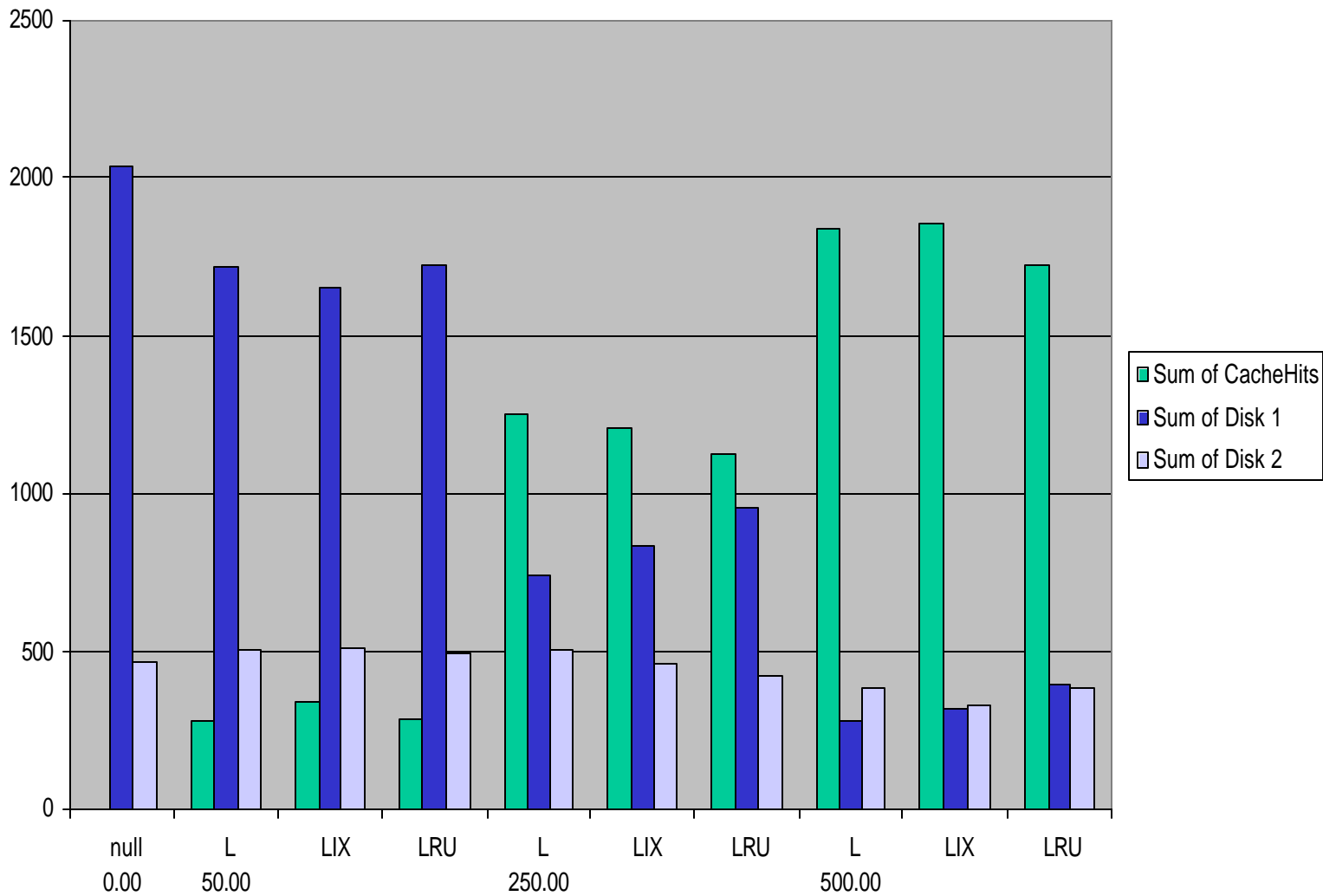
<i>ServerDBSize</i>	Number of distinct pages to be broadcast
<i>NumDisks</i>	Number of disks
<i>DiskSize_i</i>	Size of disk <i>i</i> (in pages)
Δ	Broadcast shape parameter
<i>Offset</i>	Offset from default client access
<i>Noise</i>	% workload deviation

Broadcast Disk Implementation



Broadcast Disk Implementation

<500,4500>, 0 offset, 0% noise, delta(5)



Broadcast Disk Implementation

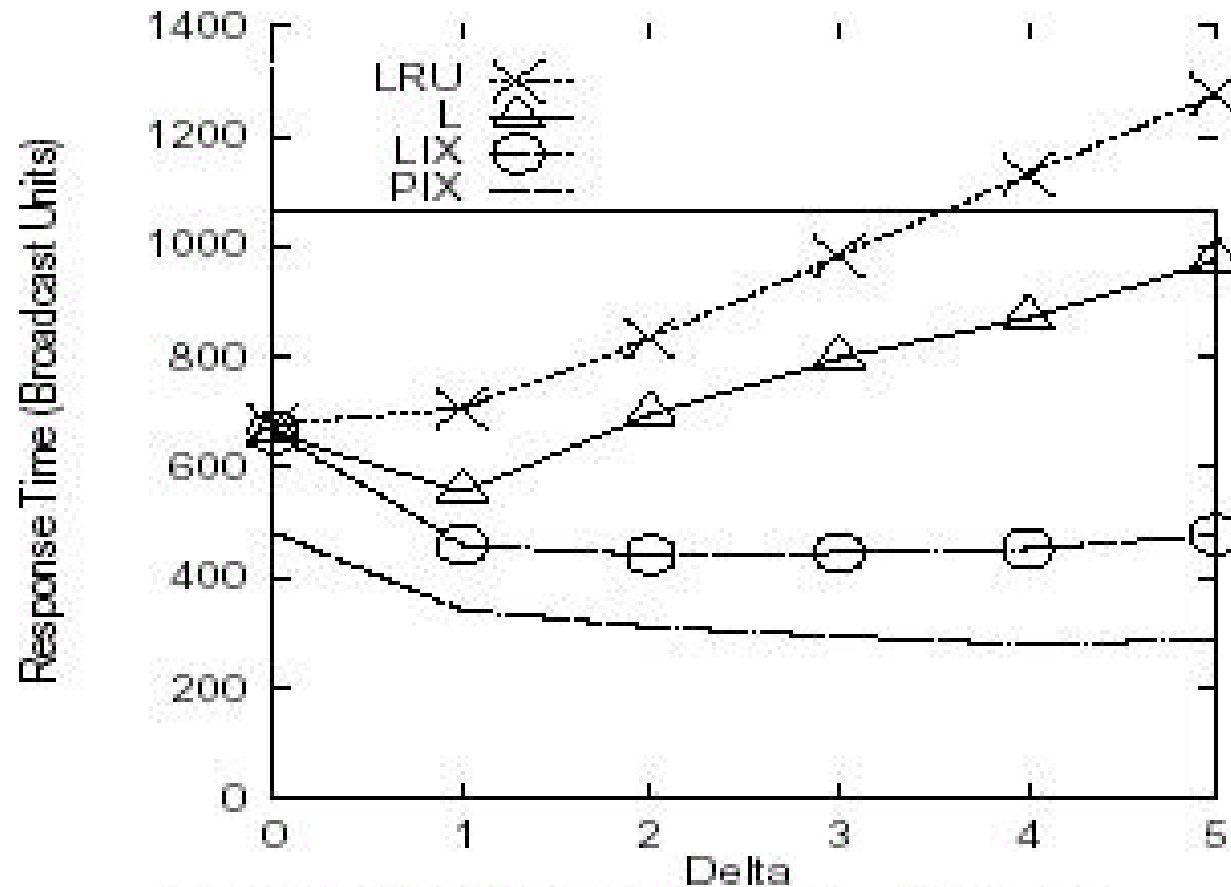
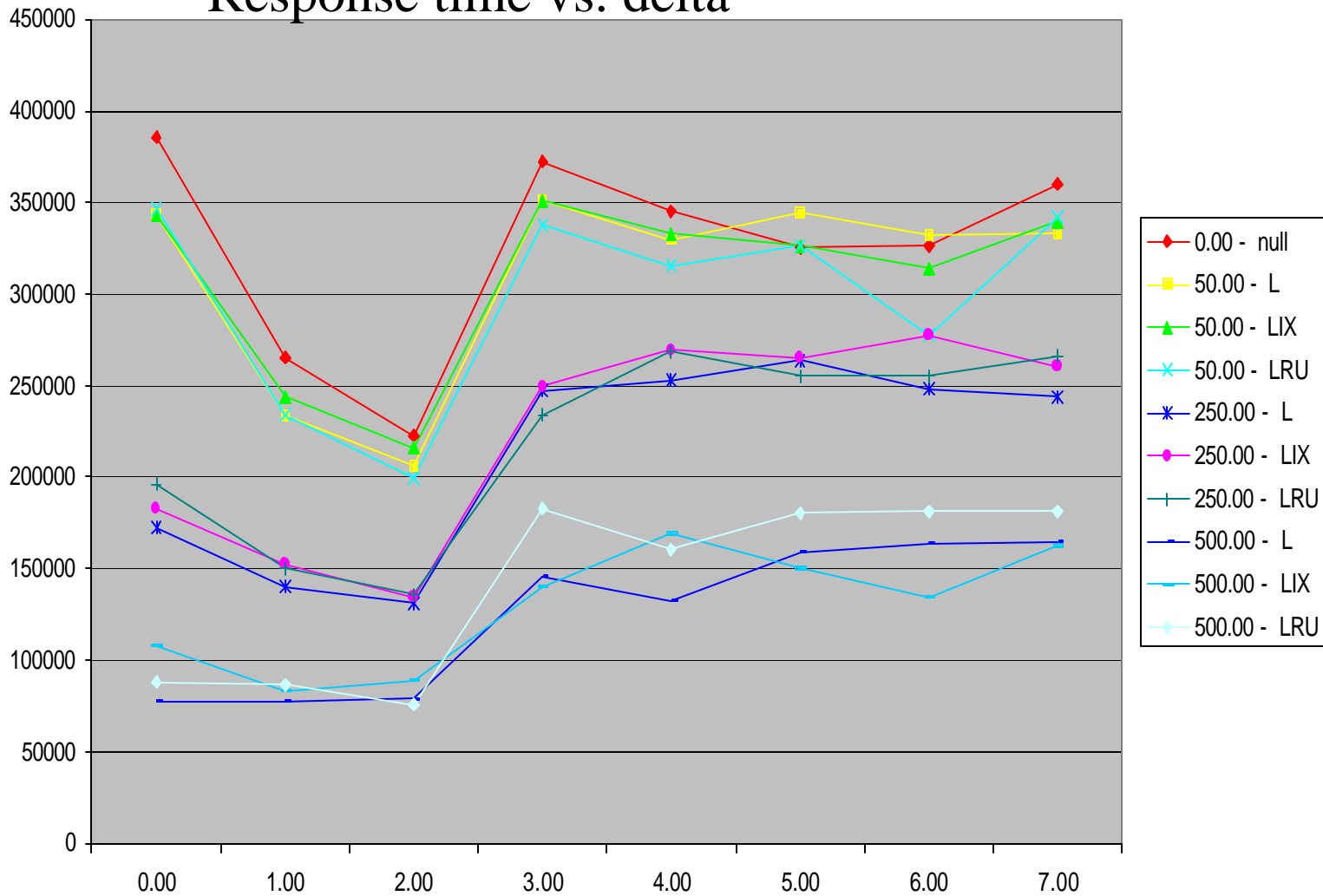


Figure 13: Sensitivity to Δ - Disk D5
CacheSize = 500. Noise = 30%

Broadcast Disk Implementation

<500,4500>, 0 offset, 0% noise

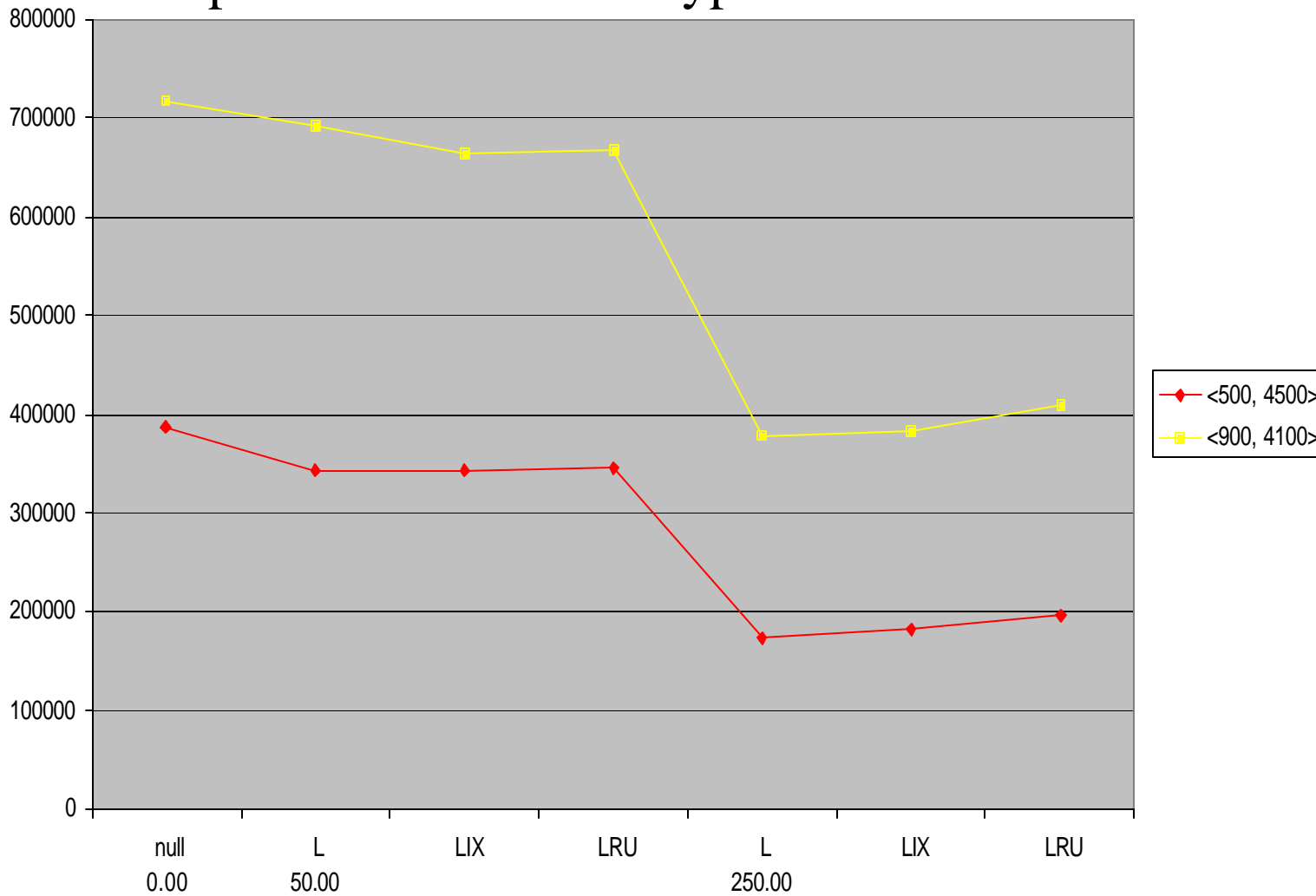
Response time vs. delta



Broadcast Disk Implementation

0 delta, 0 noise, 0% offset

Response time vs cache types/sizes.... ???????



Broadcast Disk Implementation

- Relative Frequency
 - The speed of a disk, relative to slowest disk.
 - $\text{relFreq}(i)/\text{relFreq}(N) = 1 + \alpha(N - i)$
 - N = the disk speed of the slowest disk.
- Max Chunks
 - $\text{max_chunks} = \text{LCM}(\text{relFreqs})$
- NumChunks
 - $\text{num_chunks}(i)$: the number of chunks the disk is broken into.
 - $\text{num_chunks}(i) = \text{max_chunks}/\text{rel_freq}(i)$

Broadcast Disk Implementation



Delta = 0

relFreq(1) = 1, relFreq(2) = 1

numChunks(1) = 1, numChunks(2) = 1

Mini cycle 1



Delta = 1

relFreq(1) = 2, relFreq(2) = 1

numChunks(1) = 1, numChunks(2) = 2

Mini cycle 1



Mini cycle 2



Broadcast Disk Implementation



Delta = 0

relFreq(1) = 1, relFreq(2) = 1

numChunks(1) = 1, numChunks(2) = 1

Mini cycle 1



Delta = 1

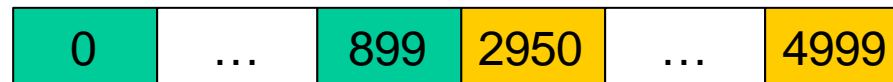
relFreq(1) = 2, relFreq(2) = 1

numChunks(1) = 1, numChunks(2) = 2

Mini cycle 1



Mini cycle 2



Broadcast Disk Implementation

- Future Work
 - Run using wireless emulator
 - Implement P and PIX
 - Let tests finish running
 - Implement Server in c or c++
 - Run on dedicated network and machines
 - Run on different packet sizes,
 - current 512 bytes
 - Implement PT and APT caching methods.

Broadcast Disk Implementation

- Conclusion
 - Time limitations and long run times, kept us from gathering all our results
 - Still L and LIX benefited from larger caches and larger deltas.
 - At small caches overhead of L and LIX not really worthwhile.