

## CORRIGENDUM

### THREE-COLOR RAMSEY NUMBERS FOR PATHS\*

ANDRÁS GYÁRFÁS, MIKLÓS RUSZINKÓ, GÁBOR N. SÁRKÖZY,  
ENDRE SZEMERÉDI

Received March 28, 2007

In February 2007 Fabricio Benevides [1] reported an easily correctable error in the proof of our main result. We wrote ([2], p. 2, lines 18–22) that one type of extremal colorings comes from an equal part blow up of a factorization of  $K_4$ . In fact, this blow up must not be necessarily equal part. A similar coloring with  $|A|, |B|, |C| \geq (1 - \alpha_1)\frac{|V(G)|}{4}$ ,  $|D| \geq \alpha_2|V(G)|$ ,  $|A| + |D| \geq (1 - \alpha_1)\frac{|V(G)|}{2}$  and coloring all edges in  $A$  with color 2 gives an extremal coloring, too. The remedy is to relax the condition in Extremal Coloring 1 (EC1) as follows, to allow one unbalanced pair  $(A, D)$ .

**Extremal Coloring 1 (with parameters  $\alpha_1, \alpha_2$ , where  $\alpha_1 \ll \alpha_2$ ).** There exists a partition  $V(G) = A \cup B \cup C \cup D$  such that

- $|A|, |B|, |C| \geq (1 - \alpha_1)\frac{|V(G)|}{4}$ ,  $|D| \geq \alpha_2|V(G)|$ ,  $|A| + |D| \geq (1 - \alpha_1)\frac{|V(G)|}{2}$ .
- The bipartite graphs  $(A \times B) \cap G_1^*$ ,  $(C \times D) \cap G_1^*$ ,  $(A \times D) \cap G_2^*$ ,  $(B \times C) \cap G_2^*$ ,  $(A \times C) \cap G_3^*$  and  $(B \times D) \cap G_3^*$  are all  $(1 - \alpha_1)$ -dense.

The proof of the fact that we can find the desired monochromatic path of length  $n$  in case we have this relaxed Extremal Coloring 1 is similar to the proof for the original EC1. For the sake of completeness we restate and prove Lemma 5 from our paper here.

**Lemma 1** ([2, Lemma 5]). *For every  $0 < \alpha_1 \ll \alpha_2 \ll 1$  there exists a positive integer  $n_0 = n_0(\alpha_1, \alpha_2)$  such that the following is true for  $n \geq n_0$ .*

---

*Mathematics Subject Classification (2000):* 05C55, 05C38

\* Combinatorica **27(1)** (2007), 35–69

If a 3-edge coloring  $(G_1, G_2, G_3)$  of  $K_{r(n)}$  is an Extremal Coloring 1 (EC1) with parameters  $\alpha_1, \alpha_2$  then there is a monochromatic path of length  $n$ .

**Proof.** First we will remove certain exceptional vertices (denote their set by  $E$ ) from the four sets  $A, B, C, D$  in EC1. A vertex  $v \in A$  is *exceptional* if one of the following is true:

$$\deg_{G_1}(v, B) < (1 - \sqrt{\alpha_1})|B|, \quad \deg_{G_2}(v, D) < (1 - \sqrt{\alpha_1})|D|,$$

or  $\deg_{G_3}(v, C) < (1 - \sqrt{\alpha_1})|C|$ .

From the density conditions in EC1 it follows that the number of these exceptional vertices is at most  $3\sqrt{\alpha_1}|A|$ . We remove these vertices from  $A$  and add them to  $E$ . Similarly, for the other three sets we define exceptional vertices and add them to  $E$ . Thus altogether (since we have at most  $2n$  vertices)

$$(1) \quad |E| \leq 24\sqrt{\alpha_1}n.$$

Next we redistribute these vertices among the 4 sets in such a way that we are not creating new exceptional vertices. Let us take the first exceptional vertex  $v$  from  $E$ , the procedure will be similar for the other vertices. Consider the  $G_1$ -neighbors of  $v$ . We may assume that these neighbors are either all in  $A \cup B$ , or in  $C \cup D$  (say they are in  $A \cup B$ ). Indeed, otherwise we can connect  $A \cup B$  with  $C \cup D$  in color  $G_1$  through  $v$  and this would give a monochromatic path in  $G_1$  of length more than  $n$  (applying Lemma 4 from [2] inside the bipartite graphs  $A \times B$  and  $C \times D$  and using  $\alpha_1 \ll \alpha_2$ ). Hence, all the edges between  $C \cup D$  and  $v$  are in colors  $G_2$  and  $G_3$ . By a similar reasoning, we may assume that  $v$  does not have  $G_2$  neighbors in both  $A \cup D$  and  $B \cup C$ , and it does not have  $G_3$  neighbors in both  $A \cup C$  and  $B \cup D$ . Thus either all the edges in  $C \times \{v\}$  are in  $G_2$ , and all the edges in  $D \times \{v\}$  are in  $G_3$ , or the other way around. Say we have the first case. Then all the edges in  $A \times \{v\}$  are in  $G_1$  and we may safely add  $v$  to  $B$ .

We repeat this procedure for all the exceptional vertices in  $E$ . Let us consider the largest set (say  $A$ ) of the four sets  $A, B, C$  and  $D$ .

**Claim 1.** If  $|B| \geq \lfloor \frac{n}{2} \rfloor$ , then there is a monochromatic path of length  $n$  in color  $G_1$  in the bipartite graph  $G_1|_{A \times B}$ .

**Proof of Claim 1.** If  $n$  is even, then take arbitrary subsets  $A' \subseteq A, B' \subseteq B$  with  $|A'| = |B'| = \frac{n}{2}$ . Applying Lemma 4 for  $G_1|_{A' \times B'}$  (the conditions of the lemma are satisfied with much room to spare) we get a monochromatic path of length  $n$  in color  $G_1$ .

If  $n$  is odd, then we must have  $|A| \geq \frac{n+1}{2}$ , since we have  $2n-1$  vertices. Then take arbitrary subsets  $A' \subseteq A, B' \subseteq B$  with  $|A'| = \frac{n+1}{2}, |B'| = \frac{n-1}{2}$ . Again

applying Lemma 4 we can find a Hamiltonian path in  $G_1|_{A' \times B'}$  beginning and ending in  $A'$ . This gives the desired monochromatic path of length  $n$  in color  $G_1$  and proves [Claim 1](#). ■

Thus we may assume that

$$(2) \quad |B|, |C|, |D| < \left\lfloor \frac{n}{2} \right\rfloor.$$

At this point we consider the colors of the edges inside  $A$ . If for the density of the  $G_1$ -edges inside  $A$  we have  $d(G_1|_A) \geq \sqrt[3]{\alpha_1}$ , then using  $\alpha_1 \ll 1$  we can clearly find a path  $P_1$  in  $G_1|_A$  that has length

$$p = \min(|A| - |B|, 2(\lceil n/2 \rceil - |B|)).$$

Remove this path from  $A$  except for one of the endpoints  $u$ . In case we have  $p < |A| - |B|$  we remove some more vertices from  $A$  until we have exactly  $|B|$  vertices left. Denote the resulting set in  $A$  by  $A'$ . Then in both cases  $|A'| = |B|$ . Again applying Lemma 4 we can find a Hamiltonian path  $P_2$  in  $G_1|_{A' \times B}$  starting with  $u$ .  $P_1$  together with  $P_2$  gives us the desired path  $P$  in  $G_1|_{A \cup B}$ . Indeed, in case  $p = 2(\lceil n/2 \rceil - |B|)$ ,  $P$  trivially has length at least  $n$ . In case  $p = |A| - |B|$ ,  $P$  is a Hamiltonian path in  $G_1|_{A \cup B}$ . By (2), in case  $n$  is even we get

$$|C| + |D| = 2n - 2 - |P| \leq 2\left(\frac{n}{2} - 1\right) = n - 2,$$

and in case  $n$  is odd we get

$$|C| + |D| = 2n - 1 - |P| \leq 2\frac{n-1}{2} = n - 1.$$

Thus in both cases

$$|P| \geq n,$$

and thus  $P$  is a monochromatic path of length at least  $n$ .

Thus we may assume  $d(G_1|_A) < \sqrt[3]{\alpha_1}$ . Similarly we may assume  $d(G_3|_A) < \sqrt[3]{\alpha_1}$ , otherwise we can find a path of length at least  $n$  in  $G_3|_{A \cup C}$ . This implies that  $d(G_2|_A) > (1 - 2\sqrt[3]{\alpha_1})$ . From this and  $\alpha_1 \ll \alpha_2$  it easily follows that the monochromatic subgraph  $G_2|_{A \cup D}$  satisfies the Pósa-condition (for nondecreasing degree sequence  $d_k \geq k+1$  for all  $k < \frac{|A \cup B|}{2}$ , see [3]) and thus has a Hamiltonian path. Similarly as above, from (2) it follows that this path has length at least  $n$ , completing the proof of the lemma. ■

The reason to relax [EC1](#) is that in Subcase 1.2 the statement about the size of the matching  $N_i$  (at least  $m_i - \sqrt{\epsilon}n$ ) is valid only if the condition

$m_i \leq |X_4| - 2|M_i|$  holds. If this condition is not true for some  $i$ , say for  $i=1$  then we easily get  $|X_4| + m_1 < (\frac{1}{2} + 4\eta)n$  which implies  $m_2 + m_3 > \frac{n}{2} - 4\eta n$ . Since we know that  $m_2, m_3 \leq \frac{n}{4} + 2\eta n$ , this implies that we have the relaxed EC1. The authors thank F. Benevides the careful reading of their manuscript.

## References

- [1] F. BENEVIDES, private communication.
- [2] A. GYÁRFÁS, M. RUSZINKÓ, G. N. SÁRKÖZY and E. SZEMERÉDI: Three-color Ramsey number for paths, *Combinatorica* **27(1)** (2007), 35–69.
- [3] L. PÓSA: A theorem concerning Hamilton lines, *Publ. Math. Inst. Hung. Acad. Sci.* **7** (1962), 225–226.

András Gyárfás

*Computer and Automation Research  
Institute  
Hungarian Academy of Sciences  
P.O. Box 63  
Budapest, H-1518  
Hungary  
[gyarf@szaki.hu](mailto:gyarf@szaki.hu)*

Gábor N. Sárközy

*Computer Science Department  
Worcester Polytechnic Institute  
Worcester, MA 01609  
USA  
[gsarkozy@cs.wpi.edu](mailto:gsarkozy@cs.wpi.edu)  
and  
Computer and Automation Research  
Institute  
Hungarian Academy of Sciences  
P.O. Box 63  
Budapest, H-1518  
Hungary*

Miklós Ruszinkó

*Computer and Automation Research  
Institute  
Hungarian Academy of Sciences  
P.O. Box 63  
Budapest, H-1518  
Hungary  
[ruszinko@szaki.hu](mailto:ruszinko@szaki.hu)*

Endre Szemerédi

*Computer Science Department  
Rutgers University  
New Brunswick, NJ 08903  
USA  
[szemer@cs.rutgers.edu](mailto:szemer@cs.rutgers.edu)*