## Lecture 6: Machine Code

• How to do Homework 2!!!!

## Homework 2

- Two parts:
  - Part 1: Use Debug to enter and run a simple machine code program
    - convert input data into 2's complement hex
    - enter data at the correct address
    - enter program at the correct address
    - run the program
  - Part 2: Write a simple machine code program, given pseudo-code
    - these instructions should be similar to those in the Part 1 problem.
    - enter and run the resulting program.

#### Part I - Example Program

Given below is a machine code program that calculates the sum of all the words in a given range of addresses in memory. The code expects that the lower bound of this range is specified in the BX register and the upper bound in the DX register. (BX holds the offset of the beginning of the data to be summed from the beginning of the data segment (DS). DX holds the offset of the last data element from the beginning of the data segment.) The sum gets stored in AX. The first 4 hex digits given on each line below represent the offset of the instruction from the beginning of the code segment. The digits after the dash are the machine code instructions. To the right are English explanations of the instructions.

0000 - 2BC0	subtract AX from itself (to make it 0)
0002 - 0307	add the word pointed to by BX to AX
0004 - 83C302	add 2 to BX (to point to the next word)
0007 - 3BD3	compare BX to DX
	(compare sets internal flags that are used by
	subsequent jump instructions)
0009 - 7DF7	if $DX \ge BX$ , then jump back to the
	instruction at 0002
000B - B8004C	this instruction and the next one return
	control to DOS
000E - CD21	

## Instruction Formats for HW2

- jump format jumping from one location in the program to another
- indirect addressing the source operand is retrieved indirectly, i.e. the operand is at the memory location pointed to by BX
- register to register format two operands, both are registers
- immediate format one operand is a constant

imm	ediate	format

#### ITR (immediate-to-register) format

For two-operand instructions in which one of the operands is a specified constant (Example: 83C302 instruction above). Its general format is:

opcode	s	w	mod	opcode	reg	immediate data
part 1				part 2		
	1	1	11			
23-18	17	16	15-14	13-11	10-8	7-0

The opcode for the add instruction in ITR format is  $100000 \ (part 1)$  and  $000 \ (part 2).$ 

The opcode for the subtract instruction in ITR format is 100000 (part 1) and 101 (part 2).

The opcode for the compare instruction in ITR format is 100000 (part 1) and 111 (part 2).

# • Lets start by looking at our program. Which instructions are immediate form?

0000 - 2BC0	subtract AX from itself (to make it 0)
0002 - 0307	add the word pointed to by BX to AX
0004 - 83C302	add 2 to BX (to point to the next word)
0007 - 3BD3	compare BX to DX
	(compare sets internal flags that are used by subsequent jump instructions)
0009 - 7DF7	if $DX \ge BX$ , then jump back to the
	instruction at 0002
000B - B8004C	this instruction and the next one return control to DOS
000E - CD21	
(actually, s look the sa	o are the last two, but those will always me so we won't worry about them now)

83C302	add 2 to BX (to point word)	to the next
10000 0 11 1	1 000 011 0000	0010
10000 s w m	od 0 0 0 r/m data	data if sw=01
→ ADD: Imr	mediate to register/m	emory
• sw = 11 - 8 bi extended to 16 • mod = 11 - r/r and specifies wi • r/m as register REG: 000 = AX 001 = CX 010 = DX 011 = BX 100 = SP r/m = 011 se • what's left? C 0000 0010 =	ts of immediate data when used) m is treated as registe hich register the inst ? Look on handout w . etc. o our register = BX Dur one byte (8 bits) = 2	(sign er field (reg) ruction uses where it says of data:

## register to register format

#### RTR (register-to-register) format

For two-operand instructions in which both the source and destination operands are registers. (Example: 2BC0 instruction above). Its general format is:

opcode	d	w	mod	dest reg	src reg
					(r/m)
	1	1	11		
15 - 10	9	8	7-6	5-3	2-0

The opcode for the **move** instruction in RTR format is 100010. The opcode for the **compare** instruction in RTR format is 001110.

The opcode for the **subtract** instruction in RTR format is 001010.

<b>.</b>	
• Lets st	art by looking at our program
again.	Which instructions are
registe	er to register form? (look for
two re	gisters and no indirection)
0000 - 2BC0	subtract AX from itself (to make it 0)
0002 - 0307	add the word pointed to by BX to AX
0004 - 83C302	add 2 to BX (to point to the next word)
0007 - 3BD3	compare BX to DX
	(compare sets internal flags that are used by
	subsequent jump instructions)
0009 - 7DF7	if $DX \ge BX$ , then jump back to the
	instruction at 0002
000B - B8004C	this instruction and the next one return
	control to DOS
000E - CD21	

## 3BD3 compare BX to DX 0011 10 11 11 01 0011 001110 dw mod reg r/m →CMP: Register/memory and register • d = 1, "to" register, d = 0, "from" register in our case, d = 1 so the reg field gives the "to" register • w = 1 - 16 bit registers (AX, BX,...not AL, BL) • mod = 11 - r/m is treated as register field (reg) • reg = 010 = DX (dest) • r/m = 011 = BX (source) (CMP does an "implied subtract" of dest ("to") - source ("from"), and sets the appropriate flags)

## jump format

#### jump format

A two-byte instruction. The first byte designates the condition on which to jump. (Example: 7D = jump if greater than or equal, in the jump instruction given in the example program. The opcode 7F means jump if greater than.) The second byte (interpreted as an 8-bit two's complement integer) gives the displacement of the jump from the *current* value of the IP. • Lets start by looking at our program again. Which instructions are in jump form? (look for the word jump)

0000 - 2BC0	subtract AX from itself (to make it 0)
0002 - 0307	add the word pointed to by BX to AX
0004 - 83C302	add 2 to BX (to point to the next word)
0007 - 3BD3	compare BX to DX
	(compare sets internal flags that are used by
	subsequent jump instructions)
0009 - 7DF7	if DX >= BX, then jump back to the
	instruction at 0002
000B - B8004C	this instruction and the next one return
	control to DOS
000E - CD21	



reason: the IP is pointing to the *next* instruction, which is at 000B. 000B - 0009 = 0002

0000 - 2BC0	subtract AX from itself (to make it 0)
0002 - 0307	add the word pointed to by BX to AX
0004 - 83C302	add 2 to BX (to point to the next word)
0007 - 3BD3	compare BX to DX
	(compare sets internal flags that are used by
	subsequent jump instructions)
0009 - 7DF7	if DX >= BX, then jump back to the
	instruction at 0002
000B - B8004C	this instruction and the next one return
	control to DOS
000E - CD21	
when executir	ng the instruction at 0009, the IP
is pointing to	the <i>next</i> instruction to be executed:
the instruction	at 000B
the instruction	
So, 000B (the	address in IP) - 0009 (the amount
to jump hoals	(which just hannons to be equal

to jump back (which just happens to be equal to the address of the current instruction) = 0002 (the address of the *next* instruction to execute.

## indirect addressing

The move instruction in your program will use indirect addressing to specify the source operand (i.e. the operand will be at the memory location pointed to by BX). The opcode for a move instruction that uses indirection is 100010; it fits into the RTR format given earlier, with mod bits = 00. (Example: 0307 instruction). • Lets start by looking at our program again. Which instructions use indirect addressing? (look for instructions where the comments say "pointed to"

0000 - 2BC0	subtract AX from itself (to make it 0)
0002 - 0307	add the word pointed to by BX to AX
0004 - 83C302	add 2 to BX (to point to the next word)
0007 - 3BD3	compare BX to DX
	(compare sets internal flags that are used by
	subsequent jump instructions)
0009 - 7DF7	if $DX \ge BX$ , then jump back to the
	instruction at 0002
000B - B8004C	this instruction and the next one return
	control to DOS
000E - CD21	

#### 0307 add the word pointed to by BX to AX

0000 0011 0000 0111

0 0 0 0 0 0 d w mod reg r/m

 $\rightarrow$ ADD: reg/memory with register to either

d = 1, reg field holds destination
w = 1 - 16 bit registers (AX, BX,...not AL, BL)
mod = 00 - there are no displacement fields in the instruction.
reg = 000 = AX (dest)
r/m = 111, EA = (BX) + Disp

EA = (BX) + Disp?
EA = effective address - the address of the word being added to AX
(BX) = contents of BX
Disp - an additional displacement field, 0 in this instruction (mod = 00)

## more on mod and r/m

- If mod = 11, then r/m is treated as a REG field. This means, you look up the r/m contents on the REG table.
- Otherwise, mod indicates if a displacement is included in the instruction.
- What's a displacement? Part of the effective address. You'll see more on displacements when we cover addressing modes.
- So mod is NOT 11, now what do we do?

### another example

#### ?? move the word pointed to by BX to DX

look for it in the instruction set list

100010d w | mod reg r/m

→MOV: reg/memory to/from register (in the data transfer section)

• d = ? well, we're moving to DX, a register. So d = 1.

•w = ? DX is a 16-bit register so w = 1

mod = ? well, we are not copying data from a register. Instead, we are copying data from a location *pointed to* by a register.
reg = ? well, we are moving the data into DX. So, look up the code for DX – 010.
r/m = ? there are a lot of choices!

#### ?? move the word pointed to by BX to DX

1 0 0 0 1 0 d w mod reg r/m

→MOV: reg/memory to/from register (in the data transfer section)

$$\begin{split} r/m &= 000, \, EA = (BX) + (SI) + DISP \\ r/m &= 001, \, EA = (BX) + (DI) + DISP \\ r/m &= 010, \, EA = (BP) + (SI) + DISP \\ r/m &= 100, \, EA = (BP) + (DI) + DISP \\ r/m &= 100, \, EA = (SI) + DISP \\ r/m &= 101, \, EA = (DI) + DISP \\ r/m &= 110, \, EA = (BP) + DISP (w/exception) \\ r/m &= 111, \, EA = (BX) + DISP \end{split}$$

So which is it? Well, lets eliminate any that use registers that are not in our instruction: SI? DI? BP? this leaves r/m = 111, EA = (BX) + DISP

#### ?? move the word pointed to by BX to DX

1 0 0 0 1 0 d w mod reg r/m

d = 1 because desination is a register w = 1 because it's a 16-bit instruction mod = 00 because there is no displacement reg = 010 for DX r/m = 111 for indirect addressing where the address is stored in BX.

100010 1 1 00 010 111 1000 1011 0001 0111 (spacing for convenience) = 8B17

## **Entering Data**

- You'll need to do the following:
  - Convert your data into hex.
     Negative numbers are represented in 2's complement.
  - Enter your data into memory at the address specified in the assignment.
  - Remember, each integer will take one word of storage (16 bits) and the bytes are stored in reverse order!

### Entering data example

- Data: 26, 14, -92
- Address for data: 1C554H (these are different from your assignment!)
- Convert the data: 26 = 001A, 14 = 000E, -92 = FFA4 (negative numbers are in 2's complement)
- Set the address: 1C554H
  - data address will an offset from DS (data segment register)
  - DS = 1C55H, offset = 4h
  - (EA = 1C550 + 4 = 1C554H)

### Entering Data (cont.)

- So, to enter the data at 1C554h
  - set DS = 1C55h
  - specify an offset of 4 when entering data in Debug (e ds:4)
  - enter each byte of data, remembering that for 16 bit values they are stored low byte, then high byte:
    - 1A 00 0E 00 A4 FF



## Entering the Program You're given the machine code for the program in part 1. You'll need to put it at the correct address.

- Address for program (*different from your homework*):
   1774Ch
- The code address will be an offset from CS (code segment register)
- CS = 1774h, offset = Ch.
- So you'll set the CS register. Then use the "e CS:C" command to enter the code.
- You'll also need to set the IP to 000Ch

EA	Code	Offset from CS
1774C	2B	000C
1774D	C0	000D
1774E	03	000E
1774F	07	000F
17750	83	0010
17751	C3	0011
17752	02	0012
17753	3B	0013
17754	D3	0014
17755	7D	0015
17756	F7	0016
17757	B8	0017
17758	00	0018
17759	4C	0019
1775A	CD	001A
1775B	21	001B
CS	1774	
IP	000C	

### From the assignment: The code expects that the lower bound of this range is specified in the BX register and the upper bound in the DX register. (BX holds the offset of the beginning of the data to be summed from the beginning of the data segment (DS). DX holds the offset of the last data element from the beginning of the data segment.) EA Data Offset from DS 1C554 1A 0004 0005

1C554	1A	0004
1C555	00	0005
1C556	0E	0006
1C557	00	0007
1C558	A4	0008
1C559	FF	0009
DS	1C55	
BX		
DX		



Offset from CS	Code		Data	Offset DS	from
000C	2B	1	·		
000D	CO		1A	C	004
000E	03		00		0005
000F	07		0E	C	0006
0010	83		00	0	0007
0011	C3		A4	0	8000
0012	02		FF	0	0009
0013	3B				
0014	D3	IP	00	0E	٦
0015	7D		00		
0016	F7		Г		7
0017	B8	BX	00	04	
0018	00	DX	00	08	٦
0019	4C			00	_
001A	CD	AX	00	00	
001B	21				_
		After:			
<sub>г</sub>		2BC0	subtract A	X from i	itself
CS	1774		(to make i	t 0)	
DS [	1C55	]			





Offset from CS	Code		Data	Offs DS	et from
000C	2B				
000D	CO		1A		0004
000E	03		00		0005
000F	07		0E		0006
0010	83		00		0007
0011	C3		A4		0008
0012	02		FF		0009
0013	3B				
0014	D3	IP	00	15	
0015	7D		0	10	
0016	F7				
0017	B8	BX	00	06	
0018	00	DV	00	08	
0019	4C	DA	00	08	
001A	CD	AX	00	1A	
001B	21				
NV - r PL – po overflo	no overflow ositive ow = sign: jump	After: 3BD3	compare 1 (sets flags	BX to I 3)	ЭХ

Offset from CS	Code		Data	Offset f DS	rom	
000C	2B		[			
000D	CO		1A	00	)04	
000E	03		00	00	)05	
000F	07		0E	00	)06	
0010	83		00	00	)07	
0011	C3		A4	00	008	
0012	02		FF	00	)09	
0013	3B					
0014	D3	IP	000	)F	]	
0015	7D	п	000		J	
0016	F7				1	
0017	B8	BX	000	06		
0018	00	DV	000	00	1	
0019	4C	DA	000	J8	]	
001A	CD	AX	001	IA	]	
001B	21		L		1	
After: 7DF7 if DX >= BX, then jump old IP = 0017, subtract 9: 000E						

Offset	~ .		-	Offse	et from
from CS	Code		Data	DS	
0000		1			
0000	2B				
000D	C0		1A		0004
000E	03		00		0005
000F	07		0E		0006
0010	83		00		0007
0011	C3		A4		0008
0012	02		FF		0009
0013	3B	1			
0014	D3	тр	00	10	
0015	7D	ш	00	10	
0016	F7				
0017	B8	BX	00	06	
0018	00	DV		00	
0019	4C	DX	00	08	
001A	CD	AX	00	28	
001B	21				
		After			
		0307	add the w	ord poi	nted to
1A + 0E			by BX to	AX	
= 28h			-		

Offset	<u>a</u> 1			Offse	t from
from CS	Code		Data	DS	
0000		1			
000C	2B				
000D	C0		1A		0004
000E	03		00		0005
000F	07		0E		0006
0010	83		00		0007
0011	C3		A4		0008
0012	02		FF		0009
0013	3B				
0014	D3	ID	00	13	
0015	7D	п	00	15	
0016	F7				
0017	B8	BX	00	08	
0018	00	DV		00	
0019	4C	DX	00	08	
001A	CD	AX	00	28	
001B	21			-	
		After:			
		83C302	add 2 to B	SX (to	
			point to th	ne next v	word)



Offset	Codo		Data	Offse	t from
from CS	Code		Data	DS	
000C	2B	1			
000D	<u>C0</u>		1A		0004
000E	03		00	(	0005
000F	07		0E	(	0006
0010	83		00	(	0007
0011	C3	1	A4	(	0008
0012	02		FF	(	0009
0013	3B				
0014	D3	IP	00	10	
0015	7D				
0016	F7				
0017	B8	BX	00	08	
0018	00	DX	00	08	
0019	4C			00	
001A	CD	AX	FF	CC	
001B	21	1		-	]
		After:			
		0307	add the w	ord poin	ited to
28h + FFA4			by BX to	AX	
= FFCCh					

Offset	<u>a</u> 1			Offs	et from
from CS	Code		Data	DS	
0000	20	1			
0000	<u>2B</u>				
	CO		1A		0004
000E	03		00		0005
000F	07		0E		0006
0010	83		00		0007
0011	C3		A4		0008
0012	02		FF		0009
0013	3B				
0014	D3	тр	00	13	
0015	7D	п	00	15	
0016	F7				
0017	B8	BX	00	0A	
0018	00				
0019	4C	DX	00	08	
001A	CD		FF	CC	
001B	21	лл	11	cc	
		After			
		83C302	add 2 to F	3X (to	
			point to th	ne next	word)
					,

Offset	Code		Data	Offse	et from
from CS	code		Dutu	DS	
000C	2B	]			
000D	C0		1A		0004
000E	03		00		0005
000F	07		0E		0006
0010	83		00		0007
0011	C3		A4		8000
0012	02		FF		0009
0013	3B				
0014	D3	IP	00	15	
0015	7D				
0016	F7				
0017	B8	BX	00	)A	
0018	00	DV		0.0	
0019	4C	DX	00	08	
001A	CD	AX	FF	CC	
001B	21				
		After:			
	a	3BD3	compare l	3X to D	X
NV – no overflow			(sets flags	)	
NG - n	egative				
overflo	w = sign:				
no jum	p!				





