## Lecture 6: Machine Code

- How to do Homework 2!!!!

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- How to
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## Homework 2

- Two parts:
- Part 1: Use Debug to enter and run a simple machine code program
- convert input data into 2's complement hex
- enter data at the correct address
- enter program at the correct address
- run the program
- Part 2: Write a simple machine code program, given pseudo-code
- these instructions should be similar to those in the Part 1 problem.
- enter and run the resulting program.

Part I - Example Program
Given below is a machine code program that calculates the sum of all the words in a given range of addresses in memory. The code expects that the lower bound of this range is specified in the BX register and the upper bound in the DX register. (BX holds the offset of the beginning of the data to be summed from the beginning of the data segment (DS). DX holds the offset of the last data element from the beginning of the data segment.) The sum gets stored in AX. The first 4 hex digits given on each line below represent the offset of the instruction from the beginning of the code segment. The digits after the dash are the machine code instructions. To the right are English explanations of the instructions.

| $0000-2 \mathrm{BC} 0$ | subtract AX from itself (to make it 0) |
| :--- | :--- |
| $0002-0307$ | add the word pointed to by BX to AX |
| $0004-83 \mathrm{C} 302$ | add 2 to BX (to point to the next word) <br> compare BX to DX <br> (compare sets internal flags that are used by |
| $0007-3 \mathrm{BD} 3$ | subsequent jump instructions) <br> if DX >= BX, then jump back to the <br> instruction at 0002 |
| $0009-7 \mathrm{DF} 7$ | this instruction and the next one return <br> control to DOS |
| $000 \mathrm{~B}-\mathrm{B} 8004 \mathrm{l}$ |  |
| $000 \mathrm{E}-\mathrm{CD} 21$ |  |

0002-0307 add the word pointed to by BX to AX
add 2 to BX (to point to the next word)
(compare sets internal flags that are used by
subsequent jump instructions)
0009-7DF7 if DX >= BX, then jump back to the
instruction at 0002
000E - CD21

## Instruction Formats for HW2

- jump format - jumping from one location in the program to another
- indirect addressing - the source operand is retrieved indirectly, i.e. the operand is at the memory location pointed to by BX
- register to register format - two operands, both are registers
- immediate format - one operand is a constant

| ITR (i <br> Fo <br> a s <br> Its | edi |  | ned <br> regist <br> instru <br> ant (E <br> $t$ is: | ate <br> format <br> ons in mple: 83 | ch $302$ | at <br> of the operand ruction above |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| opcode | S | w | mod | opcode | reg | immediate dat |
| part 1 |  |  |  | part 2 |  |  |
|  | 1 | 1 | 11 |  |  |  |
| 23-18 | 17 | 16 | 15-14 | 13-11 | 10-8 | 7-0 |
| The opcode for the add instruction in ITR format is 100000 (part 1) and 000 (part 2). <br> The opcode for the subtract instruction in ITR format is 100000 (part 1) and 101 (part 2). <br> The opcode for the compare instruction in ITR format is 100000 (part 1 ) and 111 (part 2). |  |  |  |  |  |  |



- Lets start by looking at our program. Which instructions are immediate form?

0000-2BC0 subtract AX from itself (to make it 0)
0002-0307 add the word pointed to by BX to AX
0004-83C302 add 2 to $B X$ (to point to the next word)
0007-3BD3 compare BX to DX (compare sets internal flags that are used by subsequent jump instructions)
0009-7DF7 if DX >=BX, then jump back to the instruction at 0002
000B-B8004C this instruction and the next one return control to DOS
000E - CD21
(actually, so are the last two, but those will always look the same so we won't worry about them now)

## register to register format

RTR (register-to-register) format
For two-operand instructions in which both the source and destination operands are registers. (Example: 2BC0 instruction above). Its general format is:

| opcode | d | w | mod | dest reg | src reg |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | $(\mathrm{r} / \mathrm{m})$ |
|  | 1 | 1 | 11 |  |  |
| $15-10$ | 9 | 8 | $7-6$ | $5-3$ | $2-0$ |

The opcode for the move instruction in RTR format is 100010. The opcode for the compare instruction in RTR format is 001110.

The opcode for the subtract instruction in RTR format is 001010.

- Lets start by looking at our program again. Which instructions are register to register form? (look for two registers and no indirection)

0000-2BC0 subtract AX from itself (to make it 0)
0002-0307 add the word pointed to by BX to AX
0004-83C302 add 2 to BX (to point to the next word)
0007-3BD3 compare BX to DX
(compare sets internal flags that are used by subsequent jump instructions)
0009-7DF7 if $\mathrm{DX}>=\mathrm{BX}$, then jump back to the instruction at 0002
000B-B8004C this instruction and the next one return control to DOS

000E - CD21
,

3BD3 compare BX to DX
0011101111010011
001110 dw mod reg $\mathrm{r} / \mathrm{m}$
$\rightarrow$ CMP: Register/memory and register

- $\mathrm{d}=1$, "to" register, $\mathrm{d}=0$, "from" register in our case, $\mathrm{d}=1$ so the reg field gives the "to" register
$\cdot \mathrm{w}=1-16$ bit registers (AX, BX,...not AL,
BL)
- $\bmod =11-\mathrm{r} / \mathrm{m}$ is treated as register field (reg)
- reg = $010=$ DX (dest)
$\bullet r / m=011=$ BX (source)
(CMP does an "implied subtract" of dest ("to") - source ("from"), and sets the appropriate flags)

jump format

A two-byte instruction. The first byte designates the condition on which to jump. (Example: 7D = jump if greater than or equal, in the jump instruction given in the example program. The opcode 7 F means jump if greater than.) The second byte (interpreted as an 8-bit two's complement integer) gives the displacement of the jump from the current value of the IP.

- Lets start by looking at our program again. Which instructions are in jump form? (look for the word jump)

| $0000-2 \mathrm{BC} 0$ | subtract AX from itself (to make it 0) |
| :--- | :--- |
| $0002-0307$ | add the word pointed to by BX to AX |
| $0004-83 \mathrm{C} 302$ | add 2 to BX (to point to the next word) |
| $0007-3 \mathrm{BD} 3$ | compare BX to DX <br> (compare sets internal flags that are used by <br> subsequent jump instructions) <br> if DX >= BX, then jump back to the <br> instruction at 0002 |
| $\mathbf{0 0 0 9 - 7 D F 7}$ |  this instruction and the next one return <br> control to DOS <br> $000 \mathrm{~B}-\mathrm{B} 8004 \mathrm{C}$  |
| $000 \mathrm{E}-\mathrm{CD} 21$ |  |

## 7DF7 if $D X>=B X$, then jump back to the instruction at 0002

| 01111101 | 11110111 |
| :---: | :---: |
| 01111101 | disp |

$\rightarrow$ JNL/JGE: Jump on not less/Jump greater or equal

- (we compared DX and BX in the previous instruction and the jump uses the result of the comparison)
- displacement gives the distance to add to the IP in 2's complement. In this case it's 11110111 which is a negative number!

11110111 -> $00001000+1=00001001=9$
-> jump back 9 memory locations
so where are we now? the jump instruction is at 9. But $9-9=0$, not two! reason: the IP is pointing to the next instruction, which is at 000B. 000B $-0009=0002$

## indirect addressing

The move instruction in your program will use indirect addressing to specify the source operand (i.e. the operand will be at the memory location pointed to by BX). The opcode for a move instruction that uses indirection is 100010; it fits into the RTR format given earlier, with mod bits $=00$. (Example: 0307 instruction).

0000-2BC0 subtract AX from itself (to make it 0)
0002-0307 add the word pointed to by BX to AX 0004-83C302 add 2 to BX (to point to the next word) 0007-3BD3 compare BX to DX
(compare sets internal flags that are used by subsequent jump instructions)
0009-7DF7 if $D X>=B X$, then jump back to the instruction at 0002
000B-B8004C this instruction and the next one return control to DOS
000E-CD21
when executing the instruction at 0009 , the IP is pointing to the next instruction to be executed: the instruction at 000B.

So, 000B (the address in IP) - 0009 (the amount to jump back (which just happens to be equal to the address of the current instruction) $=0002$ (the address of the next instruction to execute.

- Lets start by looking at our program again. Which instructions use indirect addressing? (look for instructions where the comments say "pointed to"

| $0000-2 \mathrm{BC} 0$ | subtract AX from itself (to make it 0) |
| :--- | :--- |
| $\mathbf{0 0 0 2}-\mathbf{0 3 0 7}$ | add the word pointed to by BX to AX |
| $0004-83 \mathrm{C} 302$ | add 2 to BX (to point to the next word) |
| $0007-3 \mathrm{BD} 3$ | compare BX to DX <br> (compare sets internal flags that are used by |
|  | subsequent jump instructions) |
| $0009-7 \mathrm{DF} 7$ | if DX >= BX, then jump back to the <br> instruction at 0002 |
| $000 \mathrm{~B}-\mathrm{B} 8004 \mathrm{C}$ | this instruction and the next one return <br> control to DOS |
| $000 \mathrm{E}-\mathrm{CD} 21$ |  |

-002 - 83
83C302
0007-3BD3
to $B$ (to point to the next word)
compare BX to DX
subsequent jump instructions)
if $D X>=B X$, then jump back to the
instruction at 0002

000E - CD21
control to DOS
-

$$
\begin{aligned}
& 0307 \text { add the word pointed to by BX to AX } \\
& 00000011 \quad 00000111 \\
& 000000 \mathrm{dw} \text { mod reg r/m } \\
& \rightarrow \text { ADD: reg/memory with register to } \\
& \text { either } \\
& \text { - } \mathrm{d}=1 \text {, reg field holds destination } \\
& \cdot \mathrm{w}=1-16 \text { bit registers (AX, BX,...not AL, } \\
& \text { BL) } \\
& \text { - } \bmod =00-\text { there are no displacement fields in } \\
& \text { the instruction. } \\
& \text { - reg }=000=\text { AX (dest) } \\
& \bullet \mathrm{r} / \mathrm{m}=111, \mathrm{EA}=(\mathrm{BX})+\mathrm{Disp} \\
& \mathrm{EA}=(\mathrm{BX})+\text { Disp? } \\
& \text { - EA }=\text { effective address }- \text { the address of the } \\
& \text { word being added to AX } \\
& \text { - }(\mathrm{BX})=\text { contents of } \mathrm{BX} \\
& \text { - Disp - an additional displacement field, } 0 \text { in } \\
& \text { this instruction }(\bmod =00)
\end{aligned}
$$

## more on mod and $\mathrm{r} / \mathrm{m}$

- If $\bmod =11$, then $\mathrm{r} / \mathrm{m}$ is treated as a REG field. This means, you look up the $\mathrm{r} / \mathrm{m}$ contents on the REG table.
- Otherwise, mod indicates if a displacement is included in the instruction.
- What's a displacement? Part of the effective address. You'll see more on displacements when we cover addressing modes.
- So mod is NOT 11, now what do we do?


## another example

?? move the word pointed to by BX to DX
look for it in the instruction set list

$$
\begin{array}{|l|l|}
\hline 100010 \mathrm{dw} & \mathrm{mod} \mathrm{reg} \mathrm{r} / \mathrm{m} \\
\hline
\end{array}
$$

$\rightarrow$ MOV: reg/memory to/from register (in the data transfer section)

- $\mathrm{d}=$ ? well, we're moving to DX, a register. So
$\mathrm{d}=1$.
- $\mathrm{w}=$ ? DX is a 16 -bit register so $\mathrm{w}=1$
- mod = ? well, we are not copying data from a register. Instead, we are copying data from a location pointed to by a register.
- reg = ? well, we are moving the data into DX. So, look up the code for DX - 010.
$\cdot \mathrm{r} / \mathrm{m}=$ ? there are a lot of choices!


## ?? move the word pointed to by BX to DX

\section*{| 100010 dw | $\bmod \mathrm{reg} \mathrm{r} / \mathrm{m}$ |
| :---: | :---: | :---: |}

$\rightarrow$ MOV: reg/memory to/from register (in the data transfer section)

$$
\begin{aligned}
& \mathrm{r} / \mathrm{m}=000, \mathrm{EA}=(\mathrm{BX})+(\mathrm{SI})+\text { DISP } \\
& \mathrm{r} / \mathrm{m}=001, \mathrm{EA}=(\mathrm{BX})+(\mathrm{DI})+\text { DISP } \\
& \mathrm{r} / \mathrm{m}=010, \mathrm{EA}=(\mathrm{BP})+(\mathrm{SI})+\text { DISP } \\
& \mathrm{r} / \mathrm{m}=011, \mathrm{EA}=(\mathrm{BP})+(\mathrm{DI})+\text { DISP } \\
& \mathrm{r} / \mathrm{m}=100, \mathrm{EA}=(\mathrm{SI})+\text { DISP } \\
& \mathrm{r} / \mathrm{m}=101, \mathrm{EA}=(\mathrm{DI})+\text { DISP } \\
& \mathrm{r} / \mathrm{m}=110, \mathrm{EA}=(\mathrm{BP})+\text { DISP (w/exception) } \\
& \mathrm{r} / \mathrm{m}=111, \mathrm{EA}=(\mathrm{BX})+\text { DISP }
\end{aligned}
$$

So which is it? Well, lets eliminate any that use registers that are not in our instruction: SI?
DI?
BP?
this leaves $\mathrm{r} / \mathrm{m}=111, \mathrm{EA}=(\mathrm{BX})+$ DISP

## ?? move the word pointed to by BX to DX

## 100010 dw mod reg r/m

$\mathrm{d}=1$ because desination is a register $\mathrm{w}=1$ because it's a 16-bit instruction $\bmod =00$ because there is no displacement reg $=010$ for DX
$\mathrm{r} / \mathrm{m}=111$ for indirect addressing where the address is stored in BX.

```
1000101100010111
1000101100010111 (spacing for convenience)
= 8B17
```


## Entering Data

- You'll need to do the following:
- Convert your data into hex.

Negative numbers are represented in 2's complement.

- Enter your data into memory at the address specified in the assignment.
- Remember, each integer will take one word of storage ( 16 bits) and the bytes are stored in reverse order!


## Entering data example

- Data: 26, 14, -92
- Address for data: 1 C 554 H
(these are different from your assignment!)
- Convert the data:
$26=001 \mathrm{~A}, 14=000 \mathrm{E},-92=\mathrm{FFA} 4$ (negative numbers are in 2 's complement)
- Set the address: 1C554H
- data address will an offset from DS (data segment register)
- DS = 1C55H, offset $=4 \mathrm{~h}$
( $\mathrm{EA}=1 \mathrm{C} 550+4=1 \mathrm{C} 554 \mathrm{H}$ )


## Entering Data (cont.)

- So, to enter the data at 1C554h
- set DS $=1 \mathrm{C} 55 \mathrm{~h}$
- specify an offset of 4 when entering data in Debug
(e ds:4)
- enter each byte of data, remembering that for 16 bit values they are stored low byte, then high byte:
- 1A 000 E 00 A 4 FF



## Entering the Program

- You're given the machine code for the program in part 1.
- You'll need to put it at the correct address.
- Address for program (different from your homework): 1774 Ch
- The code address will be an offset from CS (code segment register)
- $\mathrm{CS}=1774 \mathrm{~h}$, offset $=\mathrm{Ch}$.
- So you'll set the CS register. Then use the "e CS:C" command to enter the code.
- You'll also need to set the IP to 000Ch


## From the assignment:

The code expects that the lower bound of this range is specified in the BX register and the upper bound in the DX register. (BX holds the offset of the beginning of the data to be summed from the beginning of the data segment (DS). DX holds the offset of the last data element from the beginning of the data segment.)



| Offset from CS | Code | Data |  | set from |
| :---: | :---: | :---: | :---: | :---: |
| 000C | 2B | IP |  |  |
| 000D | C0 |  | 1A | $\begin{aligned} & 0004 \\ & 0005 \end{aligned}$ |
| 000E | 03 |  | 00 |  |
| 000F | 07 |  | 0E 0006 |  |
| 0010 | 83 |  | $00-0007$ |  |
| 0011 | C3 |  | A4 0008 |  |
| $\begin{aligned} & 0012 \\ & 0013 \end{aligned}$ | 02 |  | FF 0009 |  |
|  | 3B |  |  | 0009 |
| $\begin{aligned} & 0014 \\ & 0015 \end{aligned}$ | D3 |  | 0010 |  |
|  | 7D |  | 0010 |  |
| 0016 | F7 | BX |  |  |
| 0017 | B8 |  | 0004 |  |
| 0018 | 00 | DX | 0008 |  |
| 0019 | 4C |  |  |  |  |
| 001B | CD | AX | 001A |  |
|  | 21 | After: 0307 | add the word pointed to by BX to AX |  |
|  |  |  |  |  |  |
| CS | 1774 |  |  |  |  |
| DS | 1 C 55 |  |  |  |



| Offset from CS | Code |  | Data ${ }^{\text {O }}$ | set from |
| :---: | :---: | :---: | :---: | :---: |
| 000C | 2B |  |  |  |
| 000D | C0 |  | 1A | 0004 |
| 000E | 03 |  | 00 | 0005 |
| 000F | 07 |  | 0E | 0006 |
| 0010 | 83 |  | 00 | 0007 |
| 0011 | C3 |  | A4 | 0008 |
| 0012 | 02 |  | FF | 0009 |
| 0013 | 3B |  |  |  |
| 0014 | D3 | IP | 0013 |  |
| 0015 | 7D |  |  |  |
| 0016 | F7 |  |  |  |
| 0017 | B8 | BX | 0006 |  |
| 0018 | 00 |  |  |  |
| 0019 | 4 C | DX | 0008 |  |
| 001A | CD | AX | 001A |  |
| 001B | 21 |  |  |  |
|  |  |  |  |  |
| CS | 1774 |  | point to the n | word) |
| DS | 1 C 55 |  |  |  |


| Offset from CS | Code |  | Data Of | fset from |
| :---: | :---: | :---: | :---: | :---: |
| 000C 000D | 2B |  |  |  |
|  | C0 |  | 1A | 0004 |
| 000E | 03 |  | 00 | 0005 |
| 000F | 07 |  | 0E | 0006 |
| 0010 | 83 |  | 00 | 0007 |
| 0011 | C3 |  | A4 | 0008 |
| 0012 | 02 |  | FF | 0009 |
| 0013 | 3B | IP |  |  |
| $\begin{aligned} & 0014 \\ & 0015 \end{aligned}$ | D3 |  | 0015 |  |
|  | 7D |  |  |  |
| 0016 | F7 |  |  |  |
| 0017 | B8 | BX | 0006 |  |
| 0018 | 00 | DX | 0008 |  |
| $\begin{gathered} 0019 \\ 001 \mathrm{~A} \end{gathered}$ | 4C | DX | 0008 |  |
|  | CD | AX | 001A |  |
| 001B | 21 |  | compare BX to DX (sets flags) |  |
|  |  |  |  |  |
| PL - positive overflow = sign: jump |  | 3BD3 |  |  |








| Offset from CS | Code | IPBXDXAXAfter3BD3 | Data Ofs | set from |
| :---: | :---: | :---: | :---: | :---: |
| $\begin{aligned} & 000 \mathrm{C} \\ & 000 \mathrm{D} \end{aligned}$ | 2B |  |  |  |
|  | C0 |  | 1A | 0004 |
| 000E | 03 |  | 00 | 0005 |
| 000F 0010 | 07 |  | 0E | 0006 |
|  | 83 |  | 00 | 0007 |
| 0011 | C3 |  | A4 | 0008 |
| $0012$ | 02 |  | FF | 0009 |
| $\begin{aligned} & 0013 \\ & 0014 \end{aligned}$ | 3B |  |  |  |
|  | D3 |  | 0015 |  |
| 0015 | 7D |  |  |  |
| $\begin{aligned} & 0016 \\ & 0017 \end{aligned}$ | F7 |  |  |  |
|  | B8 |  | 000A |  |
| $\begin{aligned} & 0018 \\ & 0019 \end{aligned}$ | 00 |  |  |  |
|  | 4 C |  | 0008 |  |
| $\begin{aligned} & \text { 001A } \\ & 001 \mathrm{~B} \end{aligned}$ | CD |  | FFCC |  |
|  | 21 |  |  |  |
| NV - no overflow NG - negative overflow /= sign: no jump! |  |  | compare BX to DX (sets flags) |  |
|  |  |  |  |  |



## Result

- At the end of the program (but prior to the last two machine instructions), AX holds the result: FFCCh
- This is a 2's complement result. As we learned earlier, if a 2's complement number has a 1 in the left-most bit, it is negative.
FFCCh = 1111111111001100
$=-0000000000110011+1$
$=-0000000000110100$
$=-(2 * * 5+2 * * 4+4)$
$=-52$
- Is this correct?

$$
26+14-92=40-92=-52
$$

(if the left-most bit of your answer was zero, you would be able to simply convert it to decimal with no inverting)

## Part 2?

- Very similar to part one, except you need to figure out the machine code.
- Most of these instructions are similar to those in part 1 and can be created with minor modifications to the part 1 instructions.
- Read the assignment carefully to make sure you are putting the program and data in the correct locations!

