

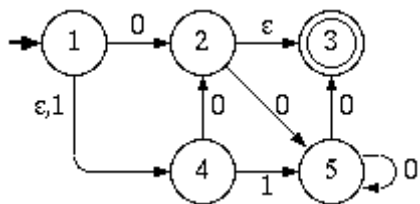
Homework #4
Please use this sheet for your answers
If you work with someone, you must write the solutions up yourself

#1. (10 Points) True or False:

- a) If $A \subseteq B \subseteq \Sigma^*$ and B is regular, then A is also regular. True False
- b) $\{w \in \{a,b\}^* \mid w \text{ contains an even number of } a\text{'s}\}$ is Regular. True False
- c) If N is an NFA, then there is an NFA accepting the complement of $\mathcal{L}(N)$. True False
- d) If $L1 \cap L2$ is regular and L1 is regular, then L2 is regular. True False
- e) If $L1 \cap L2$ is regular and L1 is finite, then L2 is regular. True False

#2. (10 Points) Show that the language P consisting of balanced parentheses is not regular. For example, $((()())())$ is a string in P.

#3. a) Minimize the following automaton and b) give the regular expression that represents the language accepted by the machine:



#4. Show that the question: “*Is regular expression r a subset of regular expression s*” is decidable.

#5. In the Product construction (See page 103 of Background Material 2 – Grahne Slides and class notes), the product transition function for machine $L \times M$ is given as:

$\delta_{L \times M} ((p,q), a) = (\delta_L(p,a), \delta_M(q,a))$ where $p \in L$ and $q \in M$. Using this and the definition of δ^* , show that $\delta_{L \times M}^*(q_L, q_M, w) = (\delta_L^*(q, w), \delta_M^*(q, w))$