

WiseMAC: An Ultra Low Power MAC Protocol for the Downlink of Infrastructure Wireless Sensor Networks

A. El-Hoiydi and J.-D. Decotignie

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Outline

- **Introduction**
- **Infrastructure Network**
- **WiseMAC**
- **ZigBee**
- **Comparison**
 - **Power-delay characteristics**
- **Conclusion**

Introduction

- **Focus on infrastructure topology**
- **Propose WiseMAC (Wireless Sensor MAC) for the downlink**
- **Trade-off power consumption and transmission delay.**
- **WiseMAC is compared to ZigBee.**

Power consumption

- **Energy efficiency is important in the sensor nodes**
- **Power consumption of transceiver in receiver mode is considerable**
- **Minimize energy waste**
 - **Idle listening – active listening to idle channel.**
 - **Overhearing – reception of a packet or part of a packet destined to another node.**

Infrastructure WSN

- **Composed of a number of access points (AP).**
- **Each access point serves a number of sensor nodes.**
- **AP is energy unconstrained**
 - **Can listen continuously**
 - **Can send any amount of signaling traffic**
 - **Exploited by WiseMAC protocol**

Traffic direction

- **Focus on low traffic situations**
- **Downlink**
 - From AP to sensor nodes
 - Transmit configuration data and query requests
 - Transmit without requiring sensor node continuously listening
- **Uplink**
 - From sensor node to AP
 - Transmit acquired data
 - AP can listen continuously with unlimited power
 - Only issue is multiple access of medium

WiseMAC

- **Medium Access Control protocol**
- **Based on CSMA with preamble sampling**
- **Sampling minimizes idle listening**
- **Exploit sensor nodes sampling schedules to minimize length of the wake-up preamble**
- **Data frames are repeated in long preambles to mitigate overhearing**

Sampling

- **Sensor nodes regularly sample the medium – listen to the radio channel for a short duration**
- **If medium found busy listen until frame is received or until idle again**
- **Sensor node sample with constant period T_w**
- **Schedule offsets are independent of each other and constant**

Preamble

- **AP transmits wake-up preamble of duration T_p in front of every data frame**
- **Ensures the receiver will be awake when the data frame arrives**
- **Provides low power consumption when channel is idle**
- **T_p is minimized by exploiting knowledge of sensor node sample schedule**

Sampling schedules

- **AP keeps an up-to-date sampling schedule of all sensor nodes**
- **Sample schedules acquired from every acknowledgment packet**
- **ACK specifies the remain time until next scheduled sampling**

WiseMAC sampling activity

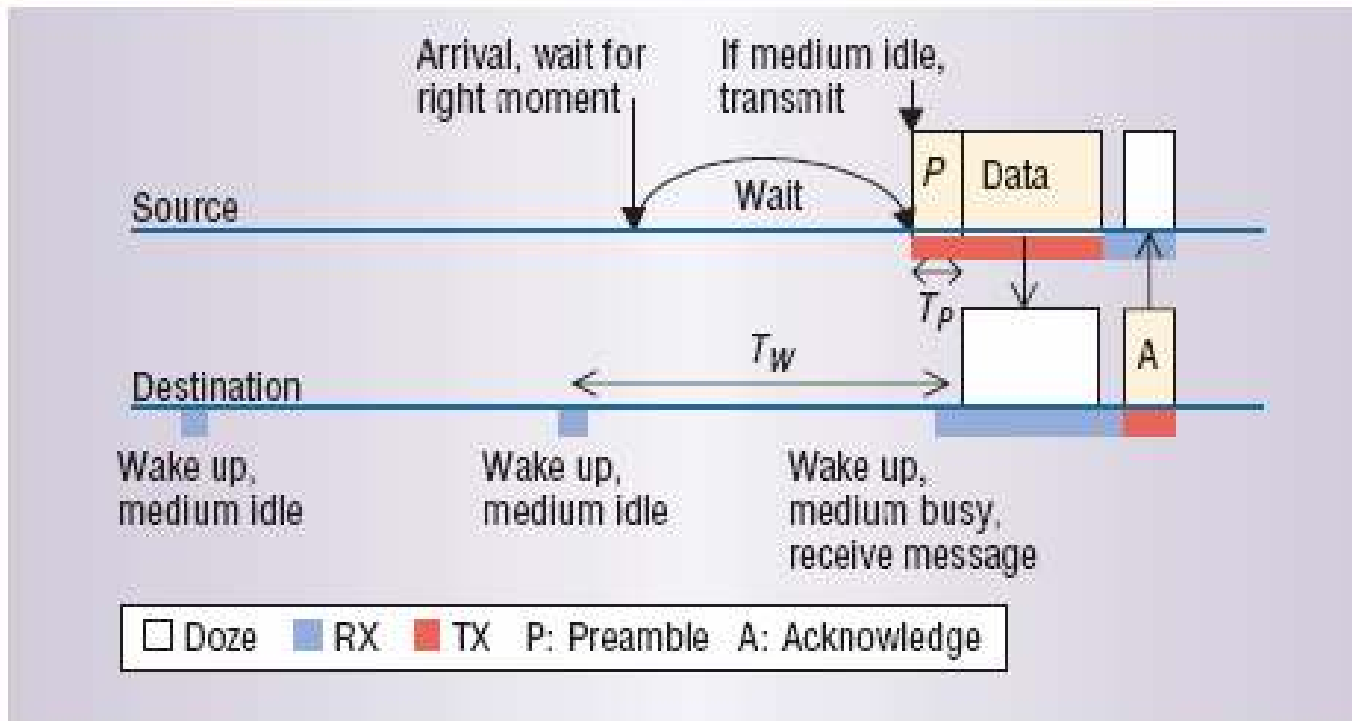


Figure 1. WiseMAC preamble minimization. A low-power, media access control protocol, WiseMAC uses a scheme that learns the sampling schedule of direct neighbors and exploits this knowledge to minimize the wake-up preamble length.

Diagram from IEEE Computer Journal feature article, WiseNET: an ultra low-power wireless sensor network solution, published by IEEE Computer Society, August 2004

Preamble duration

- T_p must compensate for drift between the clock at the AP and the sensor node
- Preamble duration must be $4\theta L$ if both quartz have a frequency tolerance of $\pm\theta$ and L is the interval between communications

Drift Compensation

- AP may be late, while node may be early, start the preamble $2\theta L$ in advance
- Because the sensor node may be late while the AP is early the duration of preamble must be $4\theta L$

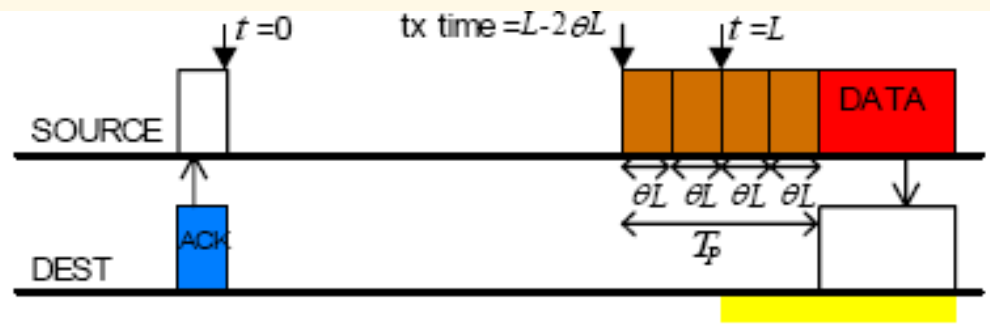


Diagram from presentation slides of Real-Time Networking Wireless Sensor Networks by Prof J.-D. Decotignie.

http://lamspeople.epfl.ch/decotignie/RTN_WSN.pdf

Drift Compensation (cont'd)

- In cases where L is very large and $4\theta L$ is larger than the sampling period T_w , the preamble length of T_w is used.

$$T_p = \min(4\theta L, T_w)$$

WiseMAC is adaptive

- In high traffic, the interval L between communications is small
- In low traffic, the interval L between communications is large, with maximum equal to T_w
- WiseMAC is adaptive to the traffic; per packet overhead decreases in high traffic conditions

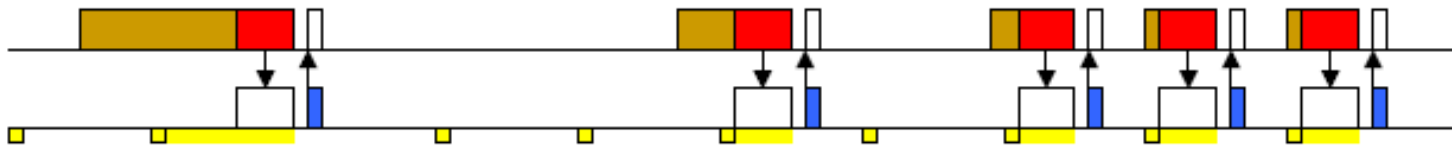


Diagram from presentation slides of Real-Time Networking Wireless Sensor Networks by Prof J.-D. Decotignie.

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High traffic conditions

- **When traffic is high overhearing is mitigated due to the preamble sampling technique and minimized preamble**
- **Short transmissions are likely to fall in between sampling instants of potential overhearers**

Low traffic conditions

- When traffic is low T_p can exceed the length of the data packet
- In which case the wake-up preamble is composed of padding bits and repetitions of the data frame

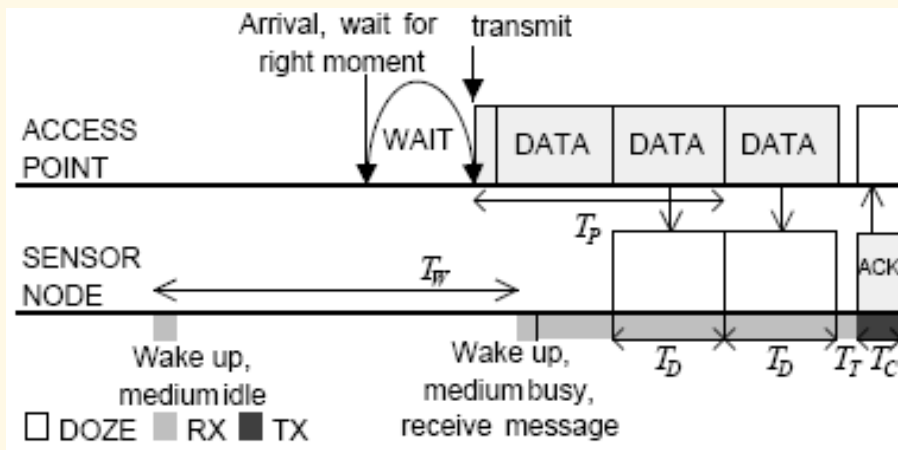


Figure 3. WiseMAC - Repetitions of the data frame in long wake-up preambles.

Frame pending bit

- ❑ In the header of the data packet
- ❑ If set, the sensor node will continue listening after having sent acknowledgment
- ❑ The AP will send the next data packet after receiving the acknowledgement
- ❑ Permits a larger wake-up interval and reduces queue delay at AP
- ❑ Cost of preamble is shared among multiple data packets

IEEE 802.15.4 ZigBee

- ❖ **WiseMAC is compared to the power save MAC protocol in ZigBee**
- ❖ **Uses central coordinator labeled access point (AP) in this document**
- ❖ **AP buffers incoming traffic**
- ❖ **AP sends periodic beacon every T_w**
- ❖ **Beacon contains address of sensor node for which data is buffered**

ZigBee Power Save Protocol

- All sensor nodes wake-up regularly to receive beacon
- Sensor node polls AP for the buffered data if the beacon contains its address
- Also uses frame pending bit in data packet header

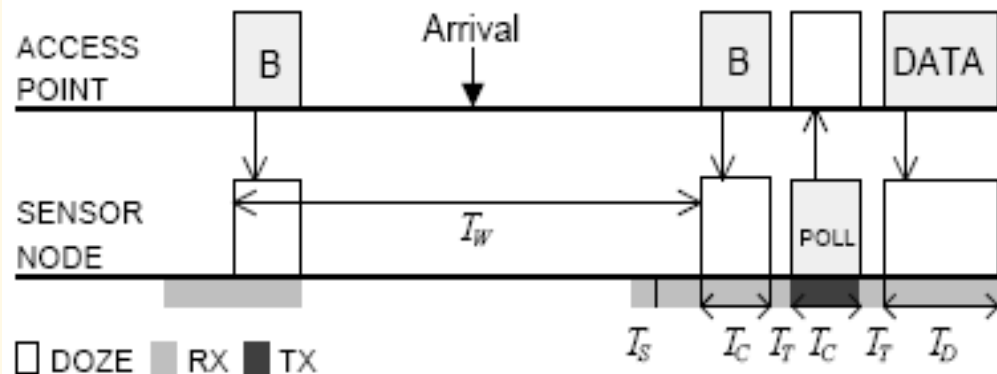


Figure 4. IEEE 802.15.4 ZigBee Power Save Protocol.

Optimize Zigbee

- ❑ For fair comparison, consider optimized version of ZigBee
- ❑ In practice polling procedure consist of POLL-ACK-DATA-ACK
- ❑ Interested in performance of basic protocol that uses beacon indication
- ❑ For low power consumption, consider POLL packet followed by DATA packet
- ❑ ACK is piggy-backed on following POLL packet

Performance Analysis

- **Model transition delays between transceiver states and power consumption in each state**
- **Transceiver states**
 - **DOZE** – The transceiver is not able to transmit nor receive, but is ready to quickly power-on into the receive or transmit state
 - **RX** – The transceiver is listening to the channel possibly receiving data
 - **TX** – The transceiver is transmitting data

Radio Model

- T_s – the setup time required to turn on the transceiver from DOZE state into the RX or TX state
- T_T – the turn-around time required to switch the transceiver between RX and TX
- P_Z, P_R, P_T – power consumed, respectively, in the DOZE, RX, and TX states
- $\hat{P}_R = P_R - P_Z$; the increment in power consumption caused by being in the RX state
- $\hat{P}_T = P_T - P_Z$; the increment in power consumption caused by being in the TX state

Traffic Model

- Population of N sensor nodes
- Downlink Poisson traffic arrives at the AP at global rate λ
- Average packet inter-arrival time at sensor node is $L = N/\lambda$
- Data packet duration is T_D
- Control packet (pollings, acks, beacons) duration is T_C
- Assume low traffic conditions

$$1/\lambda \gg T_D + T_T + T_C$$

WiseMAC Power Consumption

Average power consumed by WiseMAC

$$P_{WiseMAC} = P_Z + \frac{\hat{P}_R(T_S + 1/B)}{T_W} + \frac{\hat{P}_R(\bar{X} + T_D + T_T) + \hat{P}_T T_C}{L} + \hat{P}_R(N-1)\frac{\bar{Y}}{L}$$

Power consumed in DOZE state

Power consumed by sampling activity

Power consumed while receiving the packet and ACK it

Power consumed overhearing the packet by N-1 neighbors

where

$$\bar{X} = 2\theta L \left(1 - e^{-\frac{T_D}{4\theta L}} \right)$$

$$\bar{Y} = \frac{T_D^2 + 12T_D\theta L}{2T_W} \left(1 - e^{-\frac{T_W}{4\theta L}} \right)$$

Duration destination node listens to preamble prior to detect of start of the data frame

Average duration a potential overhearer listens to a transmission

ZigBee Power Consumption

• Average power consumed by ZigBee

$$P_{ZigBee} = P_Z + 2\theta \hat{P}_R + \frac{\hat{P}_R(T_S + T_C)}{T_W} + \frac{\hat{P}_T T_C + \hat{P}_R(T_D + 2T_T)}{L}$$

Power consumed in DOZE state

Power consumed while listening to cover the drift between AP and node

Power consumed to power on and listen to the beacon length T_c

Power consumed while polling and receiving of data packet every L seconds

Transmission delay

- The time elapsed between the arrival of a packet at the AP and the end of its transmission to the destination

$$D_{WiseMAC} = T_D + \frac{T_W}{2} \left(1 - e^{-\frac{T_W}{4\theta L}}\right) + 2\theta L \left(2 - e^{-\frac{T_D}{4\theta L}} - e^{-\frac{T_W}{4\theta L}}\right)$$

Transmission delay with WiseMAC

$$D_{ZigBee} = \frac{T_W}{2} + 2T_C + 2T_T + T_D$$

Transmission delay with ZigBee

Radio Transceiver

- Consider the transceiver used for WiseNET low power radio transceiver

Table 1. System Parameters.

Parameter	Value	Parameter	Value
P_Z	$5 \mu\text{W}$	T_D (50 bytes)	16 ms
P_R	1.8 mW	T_C (10 bytes)	3.2 ms
P_T	27 mW	N	10
T_S	0.8 ms	L	1000 s
T_T	0.4 ms	θ	30 ppm
B	25 kbps		

Power consumption and delay

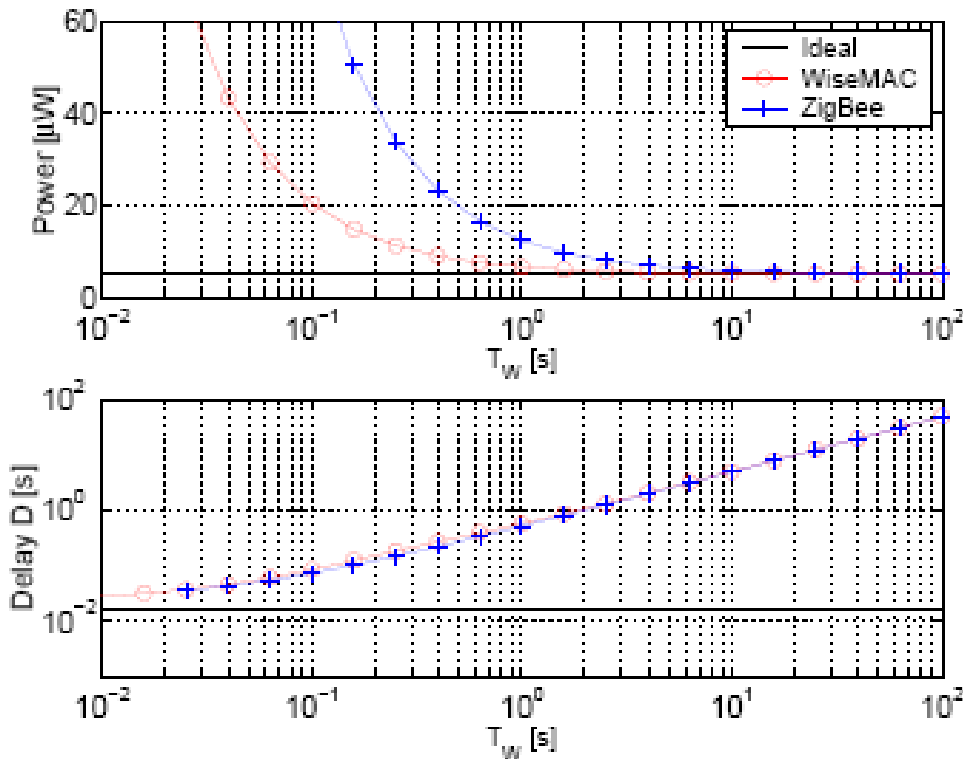


Figure 5. Power consumption and delay of WiseMAC and ZigBee as a function of the wake-up period T_w ($L = 1000$ s).

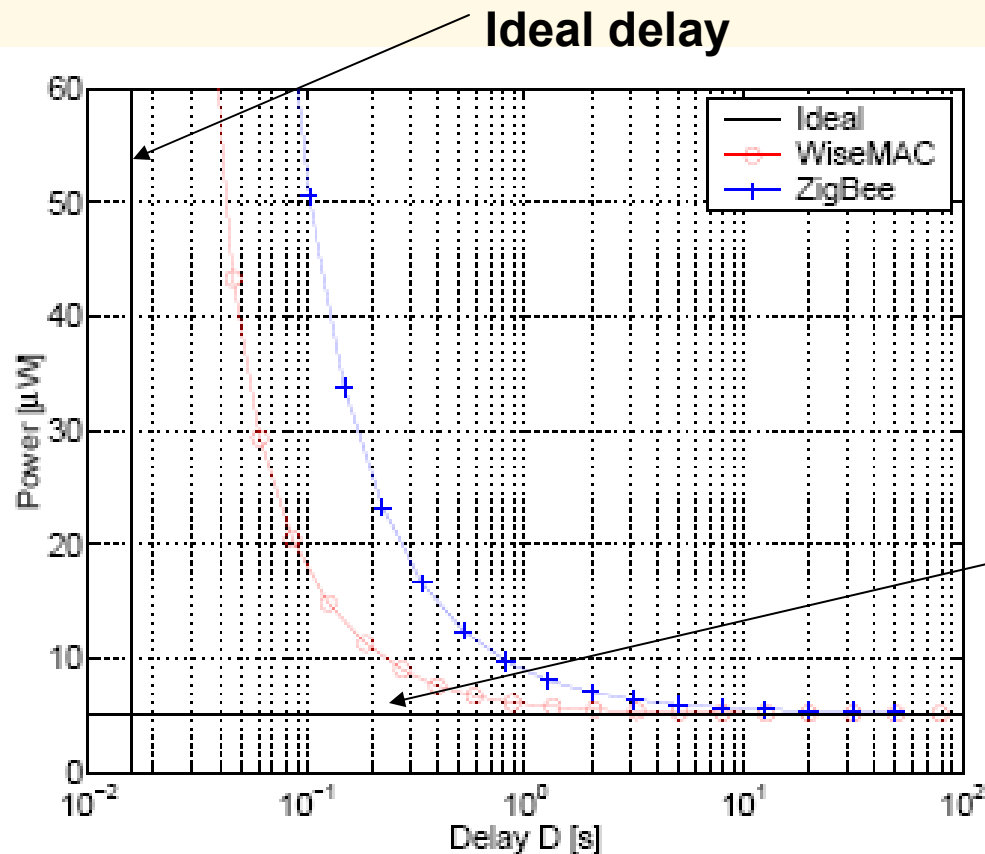
$$P_{Ideal} = P_Z + \frac{\hat{P}_R(T_S + T_D + T_T) + \hat{P}_T T_C}{L} = 5.12 \mu W$$

WiseMAC consumes less power than ZigBee

$$D_{Ideal} = 16ms$$

Trade-off between consumed power and average transmission delay

Power-delay characteristics



Combine power plot with delay plot and draw power-delay characteristics for varying T_w

Ideal power consumption

Figure 6. Power-delay characteristics of WiseMAC and ZigBee ($L = 1000$ s).

Compare wake-up schemes

- ❑ WiseMAC wake-up scheme consumes less power than the one of ZigBee
- ❑ As L approaches infinity the power consumption of WiseMAC and ZigBee becomes

$$\lim_{L \rightarrow \infty} P_{WiseMAC} = P_Z + \frac{\hat{P}_R(T_S + 1/B)}{T_W}$$

$$\lim_{L \rightarrow \infty} P_{ZigBee} = P_Z + 2\theta\hat{P}_R + \frac{\hat{P}_R(T_S + T_C)}{T_W}$$

- ❑ WiseMAC – node powers up every T_W with a duration of a radio symbol
- ❑ ZigBee – transceiver periodically receives a beacon with a duration larger than a radio symbol

Sensitivity Analysis

- **Vary the traffic and the number sensor nodes**
- **Compare WiseMAC, ZigBee, and WiseMAC***
- **WiseMAC* - a sub-optimal version where long wake-up preambles are not composed of repeated data frames**

Varying traffic

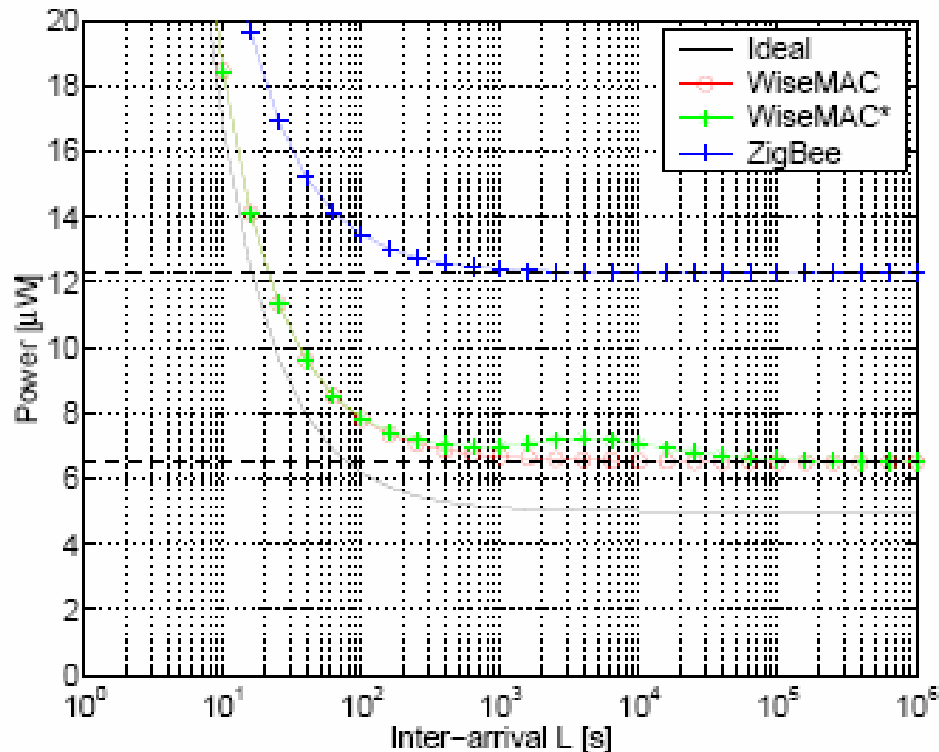


Figure 7. Power consumption of WiseMAC and ZigBee as a function of the inter-arrival L ($T_W = 1$ s).

WiseMAC has low power consumption in both high and low traffic conditions

WiseMAC* has more power consumption than WiseMAC for medium traffic – overhearing is maximized for $L \approx 4000$

Varying number of sensor nodes

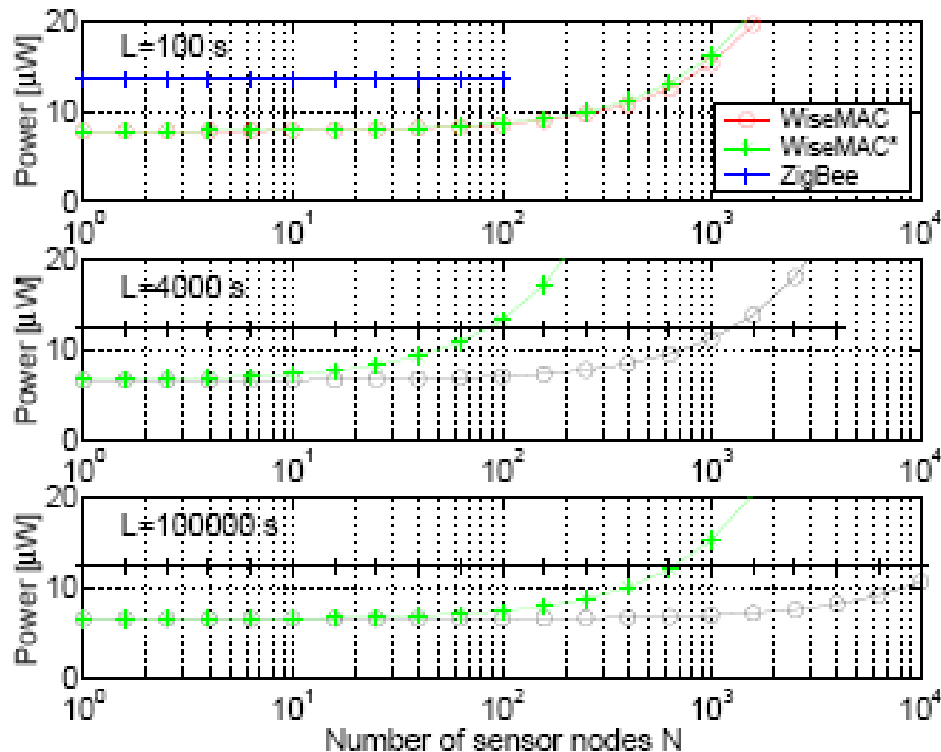


Figure 8. Scalability of WiseMAC with and without repetition of data frames in long wake-up preambles ($T_W = 1$ s).

Power consumption of ZigBee is independent of the number of nodes

Power consumption of ZigBee is independent of the number of nodes – no overhearing, scales better than WiseMAC

WiseMAC suffer from overhearing component – overhearing component is proportional to the number of nodes

Conclusion

- **Proposed WiseMAC for the downlink of infrastructure wireless sensor networks**
- **Analyzed power consumption-delay trade-off in low traffic condition and analytically compared it against ZigBee**
- **WiseMAC is more power efficient than ZigBee up to hundreds of nodes**
- **WiseMAC can provide a lower power consumption than ZigBee for the same delay**

Observations

- **Repetition of data frames in wake-up preamble explained?**