Distributed Systems

REK's adaptation of Prof. Claypool's adaptation of Tanenbaum's Distributed Systems Chapter 1

The Rise of Distributed Systems

- Computer hardware prices are falling and power increasing.
- Network connectivity is increasing.
 - Everyone is connected with fat pipes.
- It is easy to connect hardware together.
- Definition: a distributed system is
 - A collection of independent computers that appears to its users as a single coherent system.



Forms of Transparency in a Distributed System

Transparency	Description		
Access	Hide differences in data representation and how a resource is accessed		
Location	Hide where a resource is located		
Migration	Hide that a resource may move to another location		
Relocation	Hide that a resource may be moved to another location while in use		
Replication	Hide that a resource may be shared by several competitive users		
Concurrency	Hide that a resource may be shared by several competitive users		
Failure	Hide the failure and recovery of a resource		
Persistence	Hide whether a (software) resource is in memory or on disk		



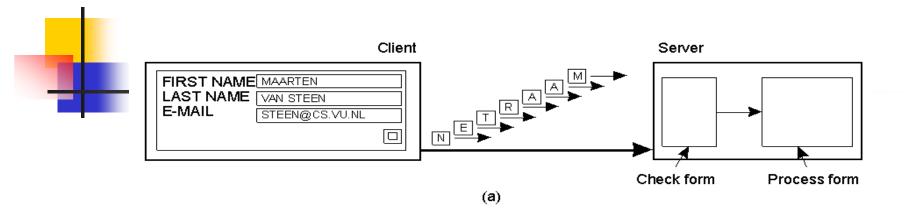
Scalability Problems

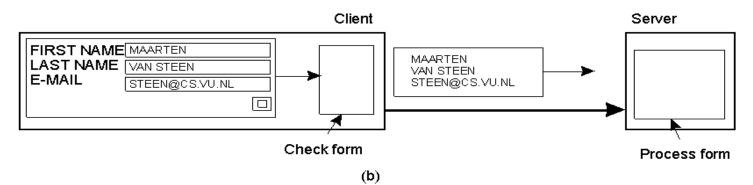
Concept	Example
Centralized services	A single server for all users
Centralized data	A single on-line telephone book
Centralized algorithms	Doing routing based on complete information

 As distributed systems grow, centralized solutions are limited.



Hiding Communication Latency

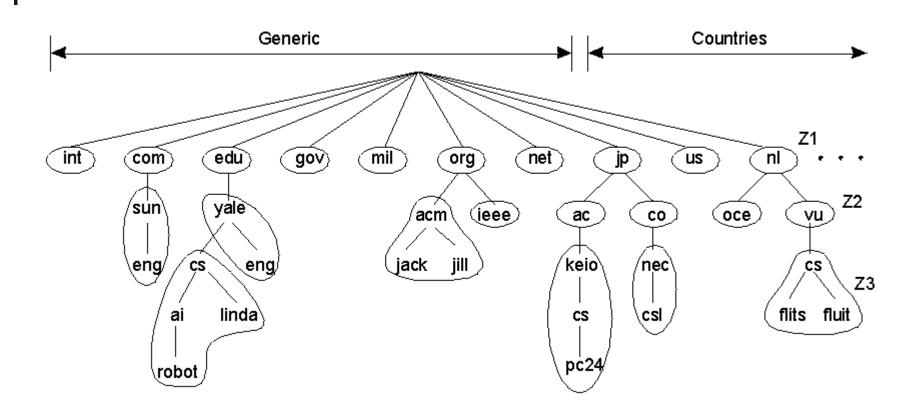




- This is especially important for interactive applications
- If possible, system can do *asynchronous communication*.
- The system can hide latencies.



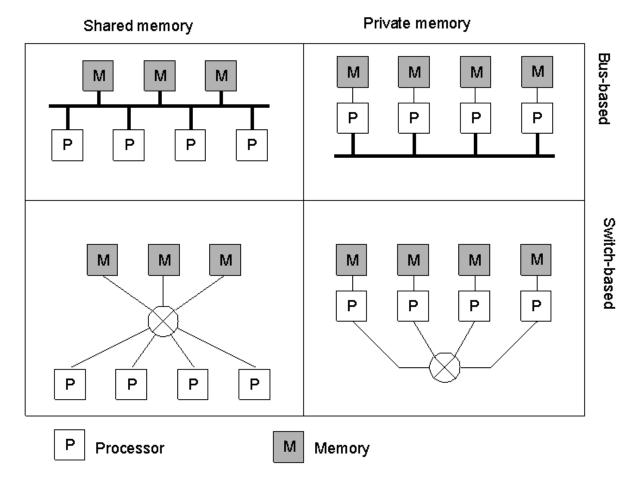
Dividing the DNS name space into zones





Hardware Concepts

Basic organizations and memories in distributed computer systems





Distributed Computing Systems

Hardware Considerations

General Classification:

- Multiprocessor a single address space among the processors
- Multicomputer each machine has its own private memory.
- OS can be developed for either type of environment.



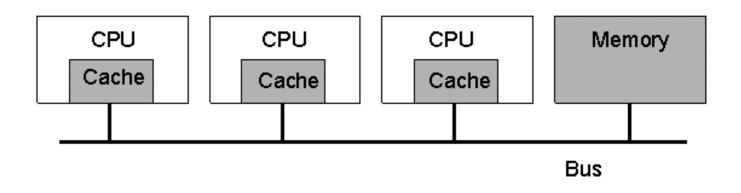
Multiprocessor Organizations

Uniform Memory Access [UMA]

- Caching is vital for reasonable performance (e.g., caches on a shared memory multiprocessor).
- Want to maintain cache coherency
 - Write-through cache :: any changes to cache are written through to memory.



Multiprocessors

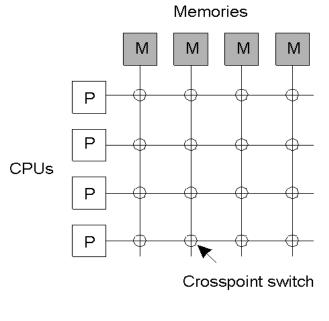


A bus-based multiprocessor.

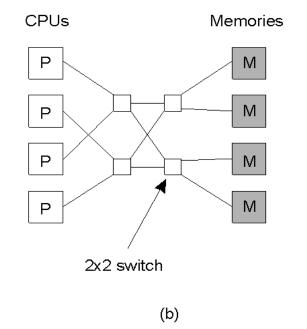


Distributed Computing Systems





(a)



A crossbar switch

An omega switching network



Multiprocessor Organizations

Non-Uniform Memory Access [NUMA]

- A hierarchy where CPUs have their own memory (<u>not</u> the same as a cache).
- Access costs to memory is non-uniform.



Replication

 Make a copy of information to increase availability and decrease centralized load.

- Example: P2P networks (Gnutella +) distribute copies uniformly or in proportion to use.
- Example: CDNs (Akamai)
- Example: Caching is a replication decision made by client.
- Issue: Consistency of replicated information
 - Example: Web Browser cache



Software Concepts

DOS (Distributed Operating Systems)

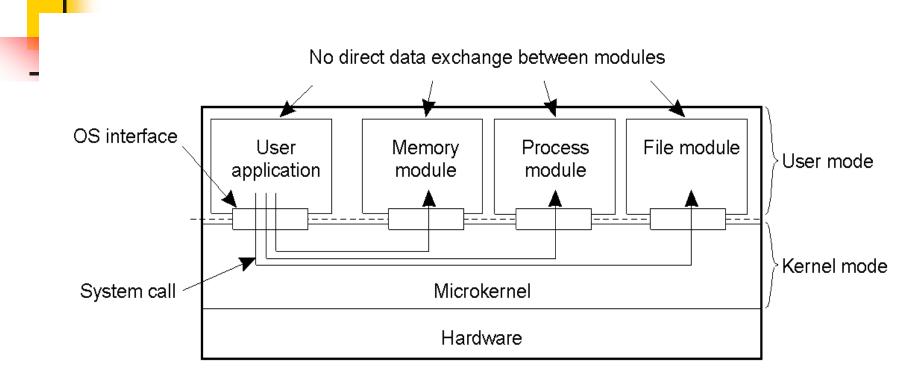
NOS (Network Operating Systems)

Middleware

System	Description	Main Goal
DOS	Tightly-coupled operating system for multi- processors and homogeneous multicomputers	Hide and manage hardware resources
NOS	Loosely-coupled operating system for heterogeneous multicomputers (LAN and WAN)	Offer local services to remote clients
Middleware	Additional layer atop of NOS implementing general-purpose services	Provide distribution transparency



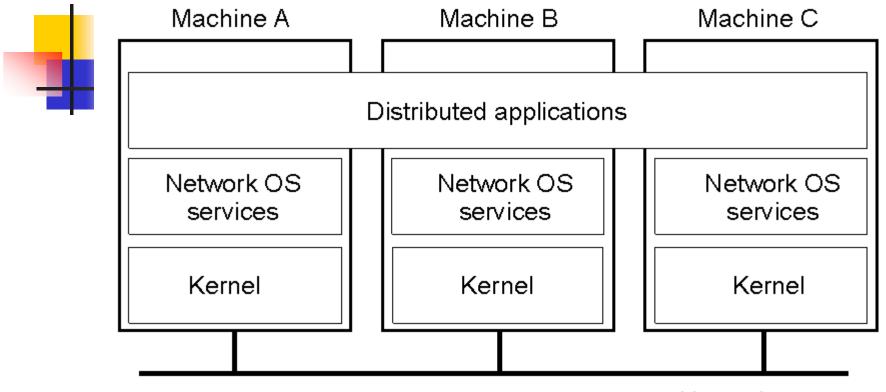
Uniprocessor Operating Systems



- Separating applications from operating system code through a microkernel
 - Can extend to multiple computers



Network Operating System

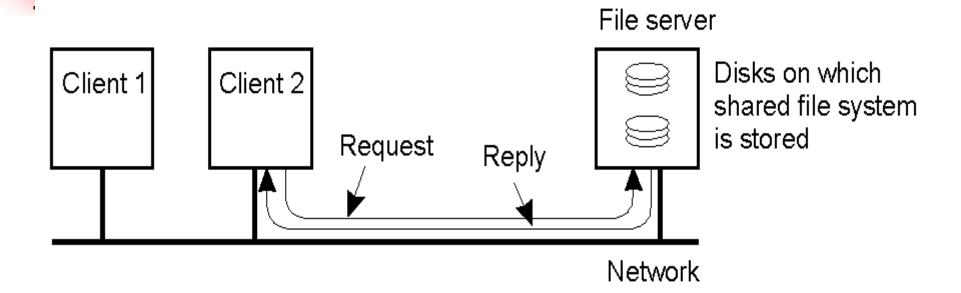


Network

- OSes can be different (Windows or Linux)
- Typical services: rlogin, rcp
 - Fairly primitive way to share files

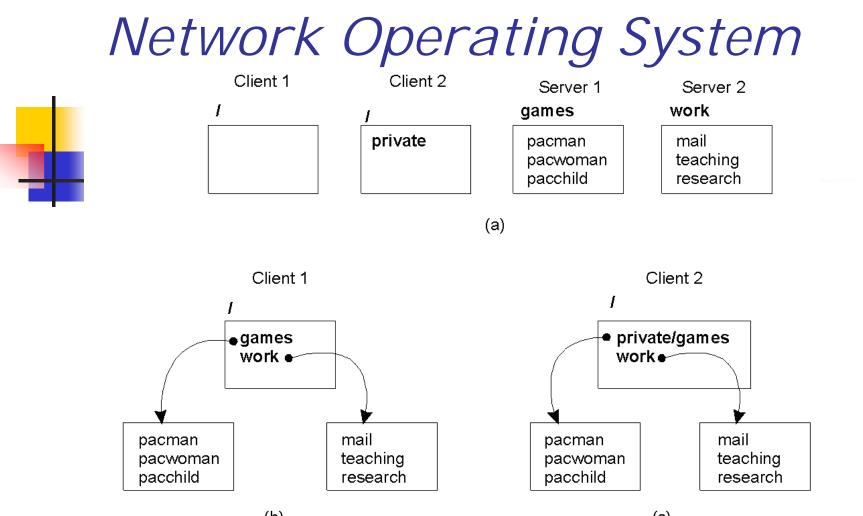


Network Operating System



- Can have one computer provide files transparently for others (NFS)
 - (try a "df" on the WPI hosts to see. Similar to a "mount network drive" in Windows)

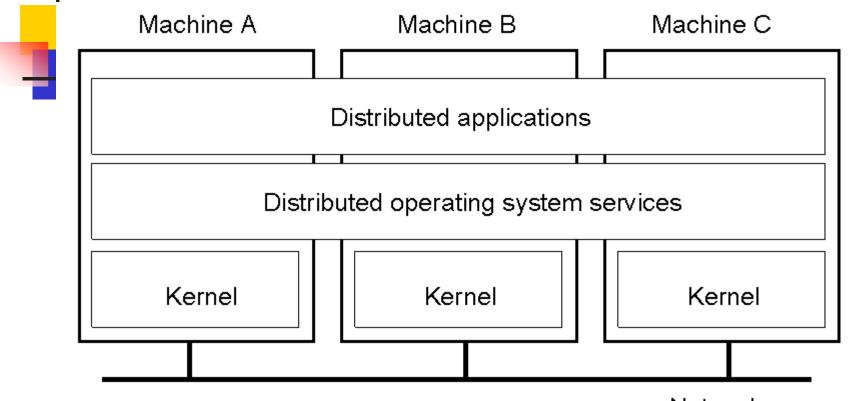




- Different clients may mount the servers in different places
- Inconsistencies in view make NOS's harder, in general for users than DOS's.
 - But easier to scale by adding computers



Distributed Operating Systems

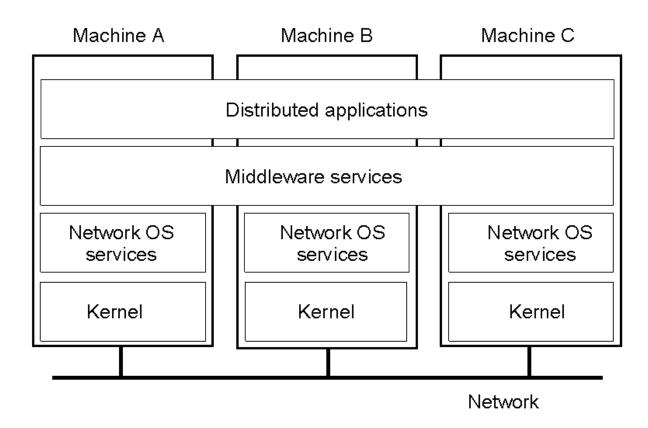


Network

- But no longer have shared memory
 - Provide message passing
 - Can try to provide distributed shared memory
 - But tough to get acceptable performance



Distributed System as Middleware





Positioning Middleware

- Network OS's are not transparent.
- Distributed OS's are not independent of computers.
- Middleware can help.

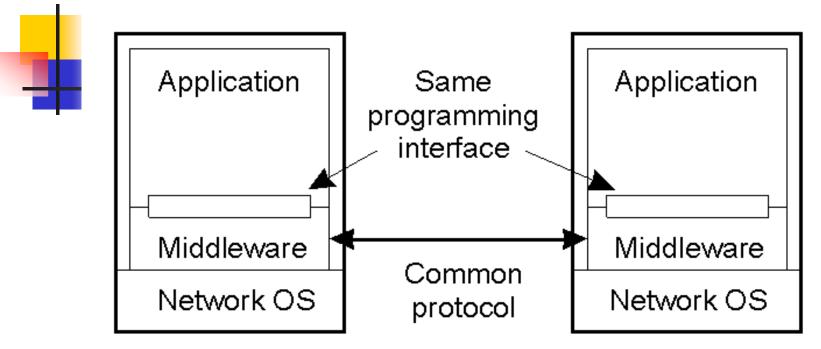


Middleware Models

- View everything as a file Plan 9.
- Less strict distributed file systems.
- Make all procedure calls appear to be local – Remote Procedure Calls (RPC).
- Distributed objects (oo model).
- The Web distributed documents.



Middleware and Openness



- In an open middleware-based distributed system, the protocols used by each middleware layer should be the same, as well as the interfaces they offer to applications.
 - If different, there will be compatibility issues
 - If incomplete, then users will build their own or use lowerlayer services (frowned upon)

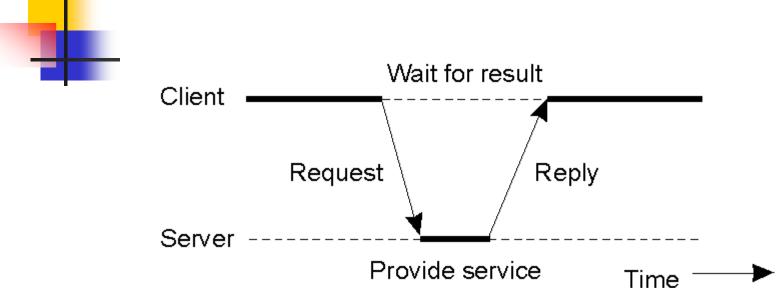


Comparison between Systems

Item	Distributed OS		Network	Middleware
Item	Multiproc.	Multicomp.	OS	-based OS
Degree of transparency	Very High	High	Low	High
Same OS on all nodes	Yes	Yes	No	No
Number of copies of OS	1	Ν	Ν	Ν
Basis for communication	Shared memory	Messages	Files	Model specific
Resource management	Global, central	Global, distributed	Per node	Per node
Scalability	No	Moderately	Yes	Varies
Openness	Closed	Closed	Open	Open



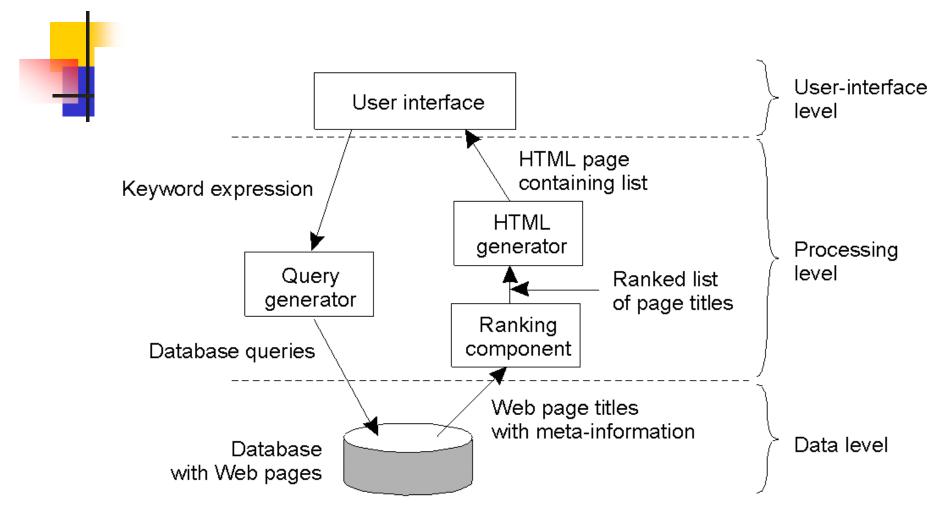
Client-Server Model



- Use TCP/IP for reliable network connection.
 - This implies the client must establish a connection before sending the first request.

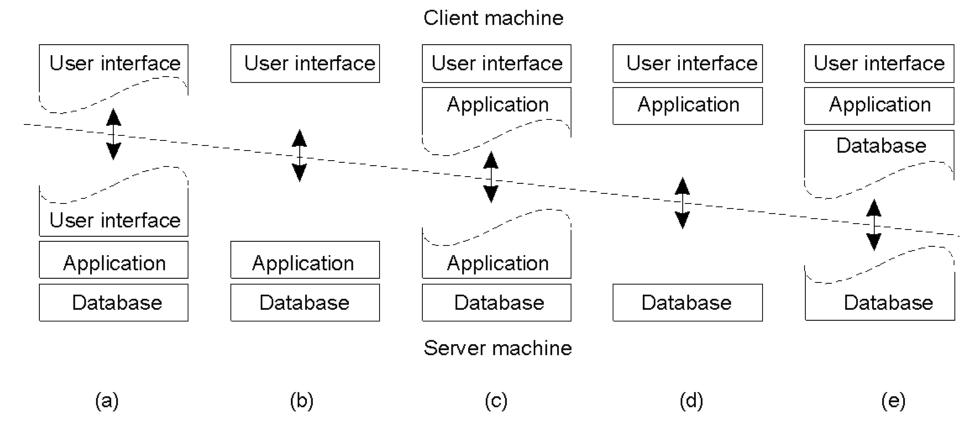


Internet Search Engine





Multitiered Architectures

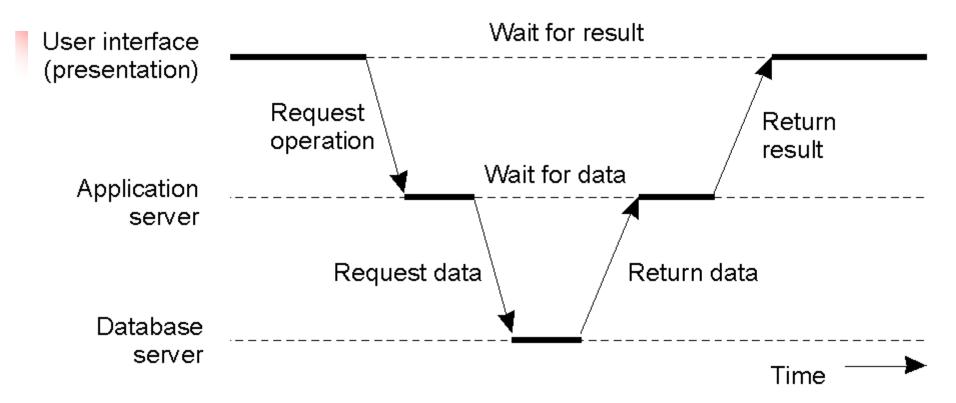


Thin client (a) to Fat client (e)

(d) and (e) popular for NOS environments



Multitiered Architectures: 3 tiers

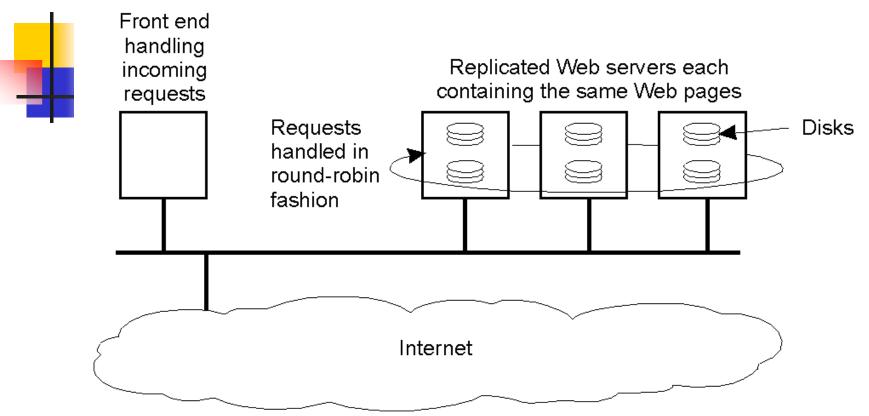


Server may act as a client

 Example would be transaction monitor across multiple databases



Horizontal Distribution



Distribute servers across nodes

• E.g., Web server "farm" for load balancing

Distribute clients in peer-to-peer systems.

